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ABSTRACT
As an instructional aid for beginning computer science courses, two systems are described which permit the automatic diagnosing and grading of student prepared problems. The first system. called SIM 610, is based on a simulator which performs by actually running student programs prepared for a simple tutorial computer used in the classroom. The simulator, which will run on any computer with even a limited FORTRAN IV capability, simulates a玉ingle address. six decimal machine with 15 basic instructions, nine index registers, and 1000 memory locations. It is capable of taking any problem and a solution prepared by the instructor and using that solution as a standard against which student problems and solutions are automatically compared and graded. The instructor can specify the weighting of factors he considers important in the grading. Diagnostic information is provided to the student on practice runs. A second system, called an Assembly Monitor, provides for the running of student machine language programs on any IBM 1130 computer. It provides a protection system against novice programers destroying resident programs and, in addition, supplies debugging aids and a grading syster very much like that for SIM 610. (JY)
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SUMMARY
As an instructional aid to beginning computer science courses, two systems are described which permit the automatic diagnosing and grading of student prepared problems. The first system is based on a simulator which performs by actually running programs prepared for a simple tutorial computer taught in the classroom. The simulator, which will run on any computer with even a limited FORTRAN IV capability, simulates a single address six decimal machine with 15 basic instructions, 9 index registers, and 1000 memory locations. Several problems which have been used in student classes are given; however, the strength in the system is that it is capable of taking any problem and it's solution, provided by the instructor, and using that solution as a standard against which student problems are automatically compared and graded. The instructor can also specify the weighting of factors he considers important in the grading. Diagnostic information is provided to the student on practice runs he makes. The system has been used over four quarters and over 600 students have run pioblems on the simulator.

A second system provides for the running of student machine language programs on any IBM 1130 computer. This system, called an Assembler Monitor, is programmed in IBM ll30 machine language itself, and will only run on that computer. It provides a protection system against novice programmers destroying resident programs in the machine and, in addition, supplies debugging aids and a grading system very much like that for the simulator.

## INTRODUCTION

In the fall of 1965 Florida Institute of Technology introduced an undergraduate degree program in Computer Science. The year 1969 saw the first graduate of this program. In addition to the more than 150 sutdents majoring in Computer Science at Florida Institute of Technology all of the 500 freshmen each year are required to take an introductory course in Computer Science which includes programming. These students pursue degree programs in Electrical Engineering, Physics, Mathematics, and Space Technology.

The Computer Science curriculum at Florida Institute of Technolngy was designed to present the technology necessary for the undergraduate student to understand computers and their usage and to become a future specialist or generalist in the field. For the non Computer Science major the introductory course taken requires the student to learn programming through actual running of programs he has prepared. For some students this is the only formal training they will receive in programming, but it provides a sufficient basis for their awn subsequent work. Others will take additional formal coursework.

Teaching of the quantities of persons taking the computer science introductory course has been a formidable problem for Florida Institute of Technology as well as at other schools in such an endeavor. Since qualified instructors are rather rare there is a natural tendency to load the good ones unmercifully in terms of the number of students they face. In such a situation the instructors find it difficult, if not impossible, to assign and evaluate a representative number of problems. Such is the motivation for a mechanized means of evaluating student problems. A mechanized system also provides for gathering and frocessing statistical data to assist the instructor in his subsequent problems assignments.

In the process of introducing the unititated to the use of electronic digital computers, and their programming in particular, a teacher or author is faced with an early decision on the specific computer he uses for illustration. He must either deal with an existing computer or develop an artificial one to demonstrate the charat teristics he deems essential. Either approach has advantages and drawbacks.

If an existing computer is taken as the illustration, a dilemma is again faced; either to choose a large machine with an extensive and sophisticated instruction repertoire, or a smaller machine with non essential characteristics imposed on it by short word length. For either case, more complexity is required than is desired to present the rudimentary concepts. The advantage of being able to demonstrate those concepts discussed on an available computer is considerable, however.

Alternatively, if an artíficial computer and its instruction repertoire are chosen as the illustrating medium, then a teaching tool can be developed exactly to the author's taste, and need only include essentials, or, it may be embellished as desired. However, the students or readers can never observe the joys of a successfully run program of their own design, or the realistic frustrations of trying to chase down a buc. The results may be like learning to drive an automobile by a correspondence course.

A compromise to the choice between a raal and an artificial computer approach is to start from the idealized artificial machine and to simulate its behavior on a real computer. In this way, programs can actually be written for the artificial computer and run (via the real one).

Work done under this contract includes the development of artificial machine language and a simulator on which it runs, and an assembler monitor system which permits ready student access to the use of an actual machine language. The simulator computer is written in FORTRAN and can be used with any computer which has a FORTRAN compiler; the assembler raonitor is for use on the IBM 1130 computer only with its machine language. The 1130 computer is in very common use in colleges and other schools and is the Floxida Institute of Technology's computer.

The automated problem set undertaken for this contract employs an artificial machine language which is simulated in the universal FORTRAN language so that programs can be written and run to demonstrate the fundamentals of machine language programming. The simulator is designated the SIM 610 (for Simulator of six decimal digit machine). Six decimal digits permits reasonable length data and instruction words. Use of decimal numbers permit the learning of concepts without the added burden of unfamiliar binary numbers, and without numerical conversions which divorce input and output numbers from internal machine
numbers and operations.
The machine language is represented in terms of an instruction set detailed in the report. The pseudo computer of the instruction set has a memory of exactly 1000 words, addressed from 000 to 999 d cimal. It has nine index registers referenced by digits 1 to 9 . It has a potential for 100 different instructions through instruction codes 00 to 99; however, only 15 of these are used. A computer word length of six decimal digits plus a sign bit (assumed + if not specified) accommodates both single address instruction and data. The 15 instructions fall into categories; data transfer, arithmetic, input/output and branching.

The Assembly Monitor System is designed to permit use of the actual machine language of the IBM-ll30 by the student in a controlled environment. This environment permits evaluation of student problems and protects the system itself from being destroyed by student program faults. Since the actual IBM-1130 machine language is rather complicated to use by an apprentice this is considered a necessary feature when assigning stadents assembly or machine language programs. Such problems are not assigned in the first introductory course which employs the SIM 610 simulator.

METHODS

## Initial Objectives

At the outset of the contract the intent witri regard to an automated set of problems was the establishment of a continuously reasable set of machine language programming problems. These problems would be of graded sophistication and difficulty and span at least two successive quarters of student experience. An evaluation and grading program was to be developed concurrently which would permit "batch" running of student programs. This program was called the Florida Institute of Technology Student Program Operating Monitor (FITSFOM).

A second task aescribed in the proposal was the preparation of a set of symbolic (assembly) language programs and a means of running these programs in batches and evaluating them also. The intenthere was to modify the IBM 1130 Assembler operating under the IBM 1130 Disc Monitor Program, a system available at many schools and colleges.

Both the evaluation programs above were to have data collecting capability on the programs run and were to perform some statistical evaluations on the results. Also both would provide feedback to the student in the form of dumps of his program.

A set of more than 60 machine language problems were developed with optimally programmed solutions and a subset of about 20 of these were picked as a set to be used in the programming courses. The problems were actually used with some of the student classes during initial work on the evaluation program and before it was ready.

A number of unanticipated difficulties arose which necessitated some revisions in the initial objectives. There are described in the following paragraphs.

A major curriculum revision occurred at Florida Institute of Technology affecting all departments and going into effect with the September 1968 term which was in the middle of the period of this grant. In this revised curriculum the courses taken by all students during the first two years are identical and it is not until the Junior year that the differences in the degree programs appear. Such a curriculum has both advantages and disadvantages for both the school and the student. From the standpoint of this grant the advantage is that not only Computer Science students, but all students at
the Institute take an introductory computer course. The disadvantage, from the grant standpoint, is that where the automated problem set was to cover a sequence of courses, it must now cover only a one quarter course and the quantity of problems which can be treated is necessarily fewer. This change did make the requirement for a mechanized handling of student programing problems mandatory for Florida Institute of Technology.

One difficulty which might have been anticipated, but was not oniginally, was that when the same problems are given to subsequent classes, the optimal solutions also pass along between the students. Thus, any finite set of problems will soon have a complete set of perfect solutions at 彐ilable within the student body so that any student who would rather copy a program than write his own finds no difficulty in doing this. This becomes particuiarly acute when the course is a mandatory one for all students and does not include just the voluntary Computer Science majors.

With the introductory programming course limited to one quarter its contents had to be very carefully evaluated so that it could best serve the needs of all students - both those Computer Science majors and the larger body some of which would not have any further formal programming. As a result it was deemed necessary to include a higher level language in the course and FORTRAN was chosen. The result is that only about half of the course is devoted to machine/symbolic language. Moreover, the machine language had to be a particularly simple one.

Student problems would really have to be prechecked before running on either the machine language or the symbolic language evaluator because they could fail to run to a finish or worse yet could destroy the evaluator or other resident programs in the computer.

New Direction to Program
As a consequence of the difficulties described, several changes occurred. A very simple machine/symbolic language was developed for an artificial but representative computer. Addressing was done in decirnal rather than binary so that concepts could be taught without the additional burden of simulateneous familiarization with another number system. Memory was limited to 1000 words.

The SIM 610 program described in this report simulates this artificial six decimal digit computer in that programs in the artificial language are executed as if the computer was real.

Instead of a formalized set of fixed problems, the approach taken was that any problem (prepared by the instructor or an advanced student, for example) could be used as a master, and the students problems would be graded against that as a standard. Thus there is no final formal set of problems; the student problems are simply made up by each instructor for each course as he needs them. Moreover, it is not assumed that the instructor's program solution is , optimal, and it is quite possible for a student grade to be higher than that of the standard provided by the instructor. Flexibility is provided for the instructor to place weighting factors on the various points to be considered in grading, changing them from problem to problem or even at different times for the same problem, depending upon where he wishes emphasis placed. For example, if he is emphasizing program running speed, a high weight can be given in the grade for fast running time as actually measured in terms of actual operations used and their execution times.

The SIM-610 simulator has been used for four quarters and with over 600 students. Surveys of student, instructor, and machine operator observations are included in this report. The Assembler Monitor System has been in informal use and aids in the writing of assembly language programs. The grading portion of the Assembler Monitor System has not been completely debugged, but since it has not had to serve large numbers of students this has not proved a problem.

## RESUITS

SIM 610 SIMULATOR SYSTEM
Philosophy for Alutomatic Grading of Student Programs
In order to grade a student's program, it is necessary to determine its operating characteristics, (i.e. what it does). It is not possible to determine what a program does except by going through it step by step, except in specialized cases. This means either running or simulating the student program. Although it would theoretically be possiole to determine other factors about a students' program not determinable simply by running or simulating it, the process involved would be too complex and time consuming to be practical.

There is one major objection to this method, however. If the student programmer makes a minor but crucial mistake anywhere in his program, his grade could be reduced to zero, even though the major part of his program works. This can be handled, however, by giving the student programmer enough debugging aids to allow him to debug his program and re-run it for a better grade. It should be noted that in practice, a computer program, no matter what methods used or how skillfully written, is worth nothing if it does not work. (We will take up the question of partially finished programs again later).

It is, therefore, necessary in order to grade a student program, to actually run it either through simulation or by allowing the execution of the instructions of the program.

If the student program is to be graded, however, the grading program must eventually regain control from the student program. This is no problem if the student's program functions properly and exits normally when finished doing the job. However, if the student's program contains an infinite (unending) loop, the grading program must be able to abort the student program and tell the student the reason for aborting. This can be best done by aborting the student program after a certain amount of run time or after a certain number of instructions have been executed (whichever is more conviently available on the system). The maximum amount of time thus set, must be large enough to allow even the inefficient student's program to complete execution; yet not allow the computer to be tied up an excessive amount of time on programs containing infinite loops. As a backup to this, it is sometimes useful to allow the operator to tell the grading program to take control. The specified method or combination of methods must be matched to the computer being used.

It should also be noted that this same instruction
count or runtime can be used later in grading the student program (see below).

It is necessary, therefore, to gain control after the student program is through executing, even if it has an infinite (endless) loop.

When the grading program has gained control, it must determine whether or not the student program has done the job assigned. In some manner the grading program must be told which problem the student is doing. It must also have been given before the student program was run, enough information to determine whether the student did the problem properly.

In order to prevent cheating, all problems should be designed so that the output is a function of the input. For example, a problem to sum the first 100 integers is not a function of an input parameter. Specificlly, the answer is a constant, 5050. The problem can be made suitable if the sum of the first "N" integers is required, where "N" is input to the student program. So long as the student does not know what value "N" will be when his program is finally graded, he must do the problem correctly in order to be assured of the correct answer.

In order to be sure that the student will not be able to cheat in this manner, the input data should be changed from practice runs before the final run of the student programs when the grades are recorded for the instructor.

In order to do the above functions, the grader must be able to feed input data to the student program. It must also have the proper answers to the problem based on this input data. The grader must also be told if some of the answers are more important than others.

What, then, should the grader do if the student programmer gets only part of the right answers? Partial credit can be given for some of the answers correct, the answers in the wrong order, or in the wrong places without too much difficulty. It should be remembered, however, that if the students are given sufficient opportunity to debug their programs, there will be little need for the grading routine to have these capabilities.

It is necessary, therefore, for the grading routine to calculate whether or not the students program did the job required on the basis of his answers being correct for the given input.

Since most students will complete a program that does the job correctly, the students grade must be based upon other factors in addition to the amount of the job completed. The best factors are those actually used to judge Fractical programs in industry: Runtime (or number of instructions executed in the student program if more easily available), and program length (ie. amount of storage space used by the student program). In addition, if the student program ended for some reason other than normal exit (ie. invalid instruction executed, excesive runtime, or other reason), then credit should be taken off.

The following formula is implemented as a weighting function to calculate the student's raw grade.

$$
G=J \times E \times(a / R+b / L+c)
$$

where
G = Raw grade to be computed;
$J=A$ factor whose value is zero if no indication was found of the job being done, and is maxium if the job was done completly correctly by the student program;
$E=A$ factor whose value is maxium if the student program ended in normal exit;
$R=$ Runtime (number of instructions executed);
L = Length of the program in core; and $a, b, c$ are positive "weighting" constants for the given problem.

One method of establishing "a", "b", and "c" is to make "a" and "b" functions of the runtime and length (respectively) of a standard program, prepared by a proficient programmer that does the job correctly. This standard program can also be used to initially calculate the proper output from the given input for use by the grader. The constant "c" provides a basis for a non vanishing grade even in the event of vanishingly snall credit for runtime, $R$, and length, $L$.

Finally, this raw grade must be curved against that of the other students doing the same problem. It is our experience that the raw grade curve can vary widely from one problem to another. Therefore only if the student's raw grade is compared to that of thers doing the same problem can his grade be curved properly. All student pro-
grams must be run for a grade before any can be given a grade in familar letter ( $A, B, C, D$, or $F$ ) or percent (100\% to $0 \%$ ) form. The raw grade (based only on the standard program for the problem) can be given each time the student program is run; even for debugging.

The grading program calculates the student's grade on the basis of whether or not he did the job, the number of instructions executed (or the runtime, if available), the length of the rrogram (how much space it uses in core), and how well his program did relative to the other students doing the same job. Moreover, the grade can be weighted by the instructor depending upon where he has placed emphasis in the programming assignment.

Finally, it is necessary to output the information thus determined by the grader. The studentis given as much information as necessary. This includes a program Iisting, reason for exit, runtime, length in core, and whether or not the program has completed the job successfully. In addition, debugging aids such as tracing all or part of the students program as it executes are included. When the programs are run for the final grade, information is supplied to the instructor so that the grades can be curved and recorded.

The SIM 610 Computer
The SIM 610 is an artikicial machine, simulated in the FORTRAN language, which will permit the student to program in machine language, and run as if his program were performing on an actual machine. The simulated computer has a word length of 6 decimal digits plus sign. When words are used for instructions, they are broken into three fields. The first two digits are the operation code, the next digit refers to any one of nine index registers, and the final three digits permit addressing any one of 1,000 addresses. Registers and data flow in the SIM 610 computer, are shown in Figure 1. Following Figure l, let us trace the operation required for the execution of a single instruction. The instruction address register will contain the address of the next instruction to be executed. Making the assumption that the tag register reads 0 (that is that none of the index registers are referenced) the address from the instruction register passes through the adder with nothing added to it and enters the memory address register. This results in the selected memory contents being placed in the memory data register, and from here it is transferred to the instruction register. While in the instruction register, the first two digits identifying the operation go to operation control to be decoded into the actual operation to be performed. The tag digit goes to the tag select switch. Here one of the index registers is identified if the tag digit is between one and nine. Finally the address is transmitted back to the memory address register through the three digit adder at which time the contents of one of the index registers may be added if it had been previously identified. The number now in the memory address register identifies the location of data in memory and this data is then brought into the memory data register. From the memory data register, the data may pass either to the input-output control, or to the transfer added and accumulator. If the operation is a print, the contents of the memory data register will actually be printed on the output print device of the real computer. If a data transfer operation is involved, such as a load accululator, the data will pass through the transfer added into the accumulator. If an arithmetic operation is involved, such as subtract from accumulator, or add to index register, the transfer adder will pass the data in the proper direction. Arithmetic operations may cause either the sign latch or the overflow latch to be set. The subsequent use of these latch indicators is described in Appendix I where each of the commands is detailed.


Fig. 1
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## Problem Definition to the Student

Each problem included in the automated problem set which students must program, must be defined to the student and to the computer simulation program so that the desired automatic evaluation can be achieved. In addition to the fundamentals of the definition, a properly solved problem must be supplied to the conputer. This solution must meet all of the specifications of the problem and should also be well programmed; that is, it should be optimum with respect to those characteristics where optimum is specified and should be near optimum in other respects. Thus, the solution should be prepared by the instructor or an otherwise well qualified programmer. This solution is called the "standard program" and all student programs are evaluated with respect to it. Nothing precludes a student bettering one or more of the parameters of the "standard program" and thus receiving a better relative score than the standard.

Characteristics which must be specified in each problem definition, provided they are appropriate are listed below:

Read: How much data must be called into the simulated computer by the program? Example: Read one card containing a number $N$ which is the order of a polynomial whose coefficients are on subsequent cards. (A total of N+2 data cards are required: 1 containing the number $N$ and $N+1$ containing the coefficients).

Store: Where are results or intermediate results to be located? Example: Calculate $f(x), f^{\prime}(x)$ and $f^{\prime \prime}(x)$ and place them in locations 200,201 and 500 respectively.

Output: What data is to be printed and in what order? Example: Print $N$ (a problem parameter) and the contents of locations 100 and 101.

Statement: A statement of the problem to be solved. Examples: (1) Read in 50 items of data and add them. (2) Print out the squares of the integers from 6 to 20 inclusive. (3) Read in $N$ numbers and sort them in increasing order of magnitude. Print out the sorted list.

Problem number: A two decimal digit number identifying the number of a problem set.

Appendix III contains some of the problems which have been assigned and solved by student classes.

Student Programs
Each student program is submitted as a deck of punched cards as follows: the first physical card in the deck is a beginning of program card, the next cards are the program proper. These are followed by an end of program card and finally by any data cards required. Format of the cards is as follows:

Beginning of program card
Column l * (asterisl)
Columns 2-7 000001
Column $8 \quad 1$ if a deck listing is desired 0 if a deck listing is not desired
Columns 9-13 five digit student number
Column 14-15 two digit problem number
Column 16 (blank)
Columns 17-51 students name (LAST FIRST)
Program card
Column $1 \quad+$ or - (blank is treated as +)
Columns 2 and 3 operation code (see Appendix I)
Column 4 tag digit (0 if no index desired)
(1-9 for index register)
Columns 5,6,7 three digit address (000-999)
End of program card
Column 1
Columns 2-7 999999
Data cards
Column 1 + or - (blank is treated as +)
Columns 2-7 six digit integer (ieading zeros if necessary).

When running programs for practice and debugging, the student should supply his own data deck following the end of data card and use an illegal problem number (e.g. 00). The data he supplies is strictly for his own use, and to satisfy himself that his program is working. If the student has suppli.ed more data cards than required, and the program finishes before using all of them, SIM simply ignores the subsequent cards as it looks for the next students beginning of program card and starts on the next program. If the student has supplied fewer data cards than required and the attempt to read another card brings out the next students beginning of program card, then the present program is terminated and the next one begun. When a program is run for credit, data cards are not supplied by the student and instead "standard" or test data is supplied by the system from disk file storage just as if it were actual cards being read on command.

The first output command executed by a student program starts a new page of printing and prints one word of data from its effective address. Execution of each subsequent output command causes one item of data to be printed on a fresh line. If the trace program is in effect, the output will be intermized with the trace, but still on a separate line.

Each run of a student program is provided with a trace of the first 25 instructions executed. Trace information (figure 2) includes on one line, the following information:

XEQNO - the number of the instruction just executed (1-25)
ADDR - the decimal address of the instruction just executed.
$C(A D D R)$ - contents of the address above (i.e. the instruction just completed.
MNEMONIC - monemonic instruction including tag and decimal address.
$\mathrm{C}(\mathrm{XR})$ - contents of index register referenced (before)
EA - effective address in instruction
$C(E A)$ - contents of effective adaress (beíore)
$C(A C C)$ - contents of accumulator (before)
$C(X R)$ - contents of index register referenced (after)
$C(E A)$ - contents of effective address (after)
SIGN - sign latch setting
OVFL - overflow latch setting
Another helpful output from a student's program run is the memory dump. This dump consists of up to 100 lines of printout, each line containing ten words (sign plus six decimal digits). Each line is headed by a decimal identifier indicating the first word of the 10 word block it contains. No blocks (lines) are printed if at least one word in the line was not changed by either writing or executing the program. Unchanged words are left blank in a line. Thus, a few lines of printout may suffice to show everything that changed in a short program. In addition (in fact prior to) the memory dump, the contents of all index registers are printed sequentially on one line. Those which were unused are again left as blank in the printout.

Additional comments which may assist the student in debugging, are provided with the trace and dump and include one of the following;

```
EXECUTION COMPLETE
PROGRAM TERMINATED DUE TO EXCESSIVE RUN TIME
INVALID INSTRUCTION ENCOUNTERED AT ----
EXECUTION TERMINATED BY INSTRUCTION AT ---- ATTEMPTING
        TO READ IST CARD OF NEXT PROGRAM INTO ----.
```

Finally, scoring information is included with calculated scores. On a grading run, the standard program weighted score is shown, otherwise it is zero.

Figure 2 is a SIM 610 diagnostic printout for the student as described in this section.


Initialization of a SIM 610 Problem
The "initialize grader" program (INITG) accepts a set of ten (10) cards containing parameters of the problem to be run, and together with other systems programs "load program" (LOADP) and "dump grader" (DUMPG) and "auxiliary initialization program" (INI2G) provides the problem description to the simulator. These ten cards and their content and function are described in the following paragraphs.

Card 0: Character set card
Columns 1 - 16: The integers and operation symbols $0123456789-b+\& * b$ where $b$ designates a blank.

Column 17: Data set Code (an integer from 1 to 6 inclusive).

The character set identifies the permissible character set and the data set designates a pair of records to be reac from "Simulator data" (SIMDT) into DATAl and DATA2 for use when the standard program executes a read card instruction.

Column 19: Final Grading Indicator. Set to lif the points and calculated grade of a student program are to be stored in SMSTU. Not used during initialization.

Cards 1 through 9 are the program description and all have the same format - 10 fields of six place integers, starting in column 1 and having two blanks between fields.

Card 1: Problem number
Field 1: Froblem number. This is the record number in the "File of standard Grades" (FSTDG).

Card 2: Read Groups
Consecutive Fields: Number of cards required in each group (NRDSR) for a number of groups up to 1 and including ten.

Card 3: Read Group Start
Consecutive Fields: The location of the first card in each read group corresponding to card 2 (LOCRD).

Card 4: Store Answers
Fields l-5: Each field gives the first of a sequence of consecutive locations in which the student program is to store answers (LCANS).

Fields 6-10: The length (number of answers) of each of the sequences starting in the respective LCANS locations above (NANSR).
Card 5: Foints Credit
Consecutive Fields: Each field stores the number of grade points credit to be given for correct answers (data matching the standard problem) for the read groups and their starting locations as given in the respective fields on cards 2 and 3. (PTCR)

Card 6: Proper answer location
Fields 1 - 5: Each field gives the number of points for placing computed answers in the proper locations (regardless of their correctness) as credit for satisfying this part of the problem specification. Proper locations are specified by the corresponding fields 1-5 and 6-10 on card 4. (PTCA)

Fields 6 - $10:$ not used.
Card 7: Correct Answers
Fields I - 5: Each field contains the number of points to be given for each correct answer found in the locations identified by card 4. (PTCC)

Fields 6 - $10:$ not used.
Card 8: Printed answer locations
Fields 1 - 5: Each field contains the number of points to be given if the correct answers are found stored in the appropriate group for printing (even if not printed in the correct sequence). (PTCW)

## Card 9:

Field l: Number of points credit if student program execute same number of card read instructions as standard program. Locations where the data read is placed is not considered here. (PTCKN)

Field 2: Number of points credit for obtaining each correct answer but storing it in an incorrect location (although within total area designated for answer storage). If an essential ingredient of the problem is intended to be sequencing or placement of results then credit points should be set to zero. (PTCO)

Field 3: Number of points credit for obtaining correct result for output but storing it in an incorrect location (although within the total area designated for output data storage). (PTWO)

Field 4: Number of answers written by standard problem. This number appears on student's dump but is not given any point value by the system. (NANS)

Field 5: The contents of this field gives the starting point within the data file for the problem under execution for the reading of simulated data cards as called for by the student (or standard) program. (FDATA)

Field 6: The number in this field establishes a maximum on the number of operations executed by a given student program. If this many steps are executed, it is assumed that the program is in a loop or is otherwise excessive in its running time and the program will be terminated. (MAXRT)

Field 7: Percent of grade for run time (steps executed for solution). (PCGRT)

Field 8: Percent of grade for program length (length of student deck). (PCGPL)

Fields 9-10: not. used.

Figure 3 ( $a$ and b) illustrate an actual set of cards from a problem set. This may be correlated with problem 3 in Appendix II.



## Operation of SIM 610

Three files must be defined in order to prepare the SIM 610 program to run a batch of student programs. These files are:

1) SIMDT This file consists of six pairs of records of 106 words each pair, and contains any simulated data required for problems to be run.
2) FSTGD This file consists of 24 records of 160 words each. Each record is associated with one problem; thus 24 different problems may be evaluated in one batch.
3) SMSTU This file consists of 800 records of 40 words each. Each record is associated wi.th one student; thus 800 student's programs may be evaluated in one batch (assuming each student has only one program.)

An initial_zation is required at the outset of a batch run in order to: a) assure that grades associated with any problem numbers undefined for the current batch give a zero grade (if not initialized, a meaningless result could occur when an undefined problem number was attempted) and b) set the "pointer" in the first record of SMSTU to the first student record (as each student deck is processed, data on his program are placed in the SMSTU at the next available position. The pointer keeps track of the next available position so that batching of student problems for grading may actually take place over more than one time on the computer.)

Ioading of data into one of the records of the SIMDT is done by the INDFG subprogram. A character set card, with the symbols
$0123456789-b+\& * b \quad(b$ is $a$ blank space)
in columns 1 through 16 and a digit $l$ through 5 in column 17 to designate which pair of records is to be loaded, must precede the data cards. This is followed by up to 106 pairs of data cards which will be entered into the designated records.

Now the SIM system is ready to initialize or run problems. For initialization, INITG is executed, and reads for each problem a character set card, nine problem definition cards, a data set of 106 cards if column 17 of the character set card was zero (did not indicate one of the six prestored DATA sets in SIMDT), and a standard program including beginning and end of program cards. Further details are given under Initialization. For running a series of students programs, STRTG is executed, which requires only one character definition card, followed by a DATA set deck if the character card so indicates, and then the student programs
stacked one after another. Normally students are given some time to debug their programs, and the results may not be desired to be recorded beyond the listing - dump which is given to the student. This will occur if anything but a 1 is in column 19 of the character set card. When the deadine for finished student programs has passed, STRTG is run using a 1 in column 19, and the student's student number, name, points received (3 categories) and raw grade are stored on a record of SMSTU for each program, except those with mispunched cards (such as a number in column l), which are not executed or dumped.

Whereas initialization may be terminated at the end of the present program by turning off sense switch 2, no provision is presently made for exiting form SIM 610 in normal operation (under STRTG), since runs are generally of long duration and abnormally terminating a FORTRAN program is simple with most computers.

AEter a class or group of students programs have been run for grading, the file SMSTU should be dumped to cards for reduction to instructor-useable forms. The program DMPFG accomplished this, and also makes a listing. This gives the programs in the order run, and is useful for finding decks or listings (if not yet returned/given to instructor) or identifying mispunched programs, which are not run, and are in the deck but not on the list. The deck is used in conjunction with a simple listing program and a card sorter, as described below.

The cards may now be sorted in ascending raw grade order and separated by problem, giving a list useful for marking grade divisions; they may be resorted in alphabetical order or student number order for instructor!s convenience. An advantage over on-disk sorts arises if correction is desired of cards which do not have last name first, or have other obtainable data missing. If more than one class is represented, the cards may be sorted on the field(s) chosen to distinguish classes, and each class deck listed in various sequences to the taste of the professor. In fact, the separate decks may be given to each instructor to cross index as he wills. Note that the original file is still available until INTFG is used to clear it. (Caution: if same problems are to be reused, references to the file FSTDG should be removed from INTFG, or else this will also be cleared; however, it is normally desired to change to a new data set both before and after grading, thus requiring reinitialization of the problems anyway).

## ASSEMBLER MONITOR SYSTEM

The Assembler Monitor System differs from the SIM 610 system in a number of ways. First of all, the Assembler Monitor System uses an actual computer language -- that of the IBM 1130, a small general purpose computer, and thus can only be used on an IBM 1130. As was discussed before, the SIM 610 system can be used on any computer that has a FORTRAN compiler.

There are some advantages to the Assembler Monitor System (AMS), however, Unlike SIM, AMS can use subroutines, including all subroutines available for the system. This also means that AMS allows more flexibility in input/output and allows for problems of much greater complexity.

Two further uses for AMS were found during development. Iike most small computers, the IBM 1130 has no memory protect hardware, and no available software to provide this feature. Therefore, we developed as part of AMS a software memory protect to prevent the student from accidentally destroying the Assembler Monitor itself, or the core-resident portion of the IBM supplied monitor-supervisor system. This feature of AMS has proved useful in itself as a debugging aid for the IBM ll30, especially for hard-to-debug assembler language programs and subroutines. Secondly, it was found that AMS could monitor FORTRAN programs on the IBM 1130 just as easily as assembler language programs, thus opening the way to additional uses for the system.

This memory protect software, a necessary part of the Assembler Monitor System, is an extremely complex system in itself. It comprises most $\dot{f}$ the $A M$ program, which consists of more than 1000 cards. It is. written in the assembler language of the IBM 1130 .

The portion of this report on the Assembler Monitor System is presented in the form of descriptive handouts to those using the system, and has worked quite effectively. Each of the subsequent sections is such a handout.

Calling the Assembler Monitor System
In order to put your program under the control of the Assembler Monitor system, it is only necessary to call AM from your program, giving it the problem number and your student number. However, there are two pitfalls that must be avoided:

1. Your call to $A M$ must be physically the first CALL or IIBF in your mainline program.
2. Your call to $A M$ should be the first executed statement in your mainline program. Should any instructions be executed before the call to AM, they will not be under control of the Assembler Monitor System.

Your input and output are in COMMON, located at the very end of core. In order to set aside this space at the end of core, you must use a COMMON statement.

You should not attempt to call AM more than once in any given program. An attempt to do so will result in the Assembler Monitor System suppressing further execution.

Below is a sample program including ca:ling sequence for the Assembler Monitor System. (Numbers in next line are card column positions.)

| 1 | 21 | 27 | 35 | 42 |
| :--- | :--- | :--- | :--- | :--- |

$/ / J O B$
//ASM
*COMMON 48
*LIST
(Note: 48 is a sample number, only) (Note: optional)

Program, constants, etc.; not including CALL or LIBF statements.

| START | CALI | AM |
| ---: | :--- | :--- |
| DC | PRNO |  |
| DC | STNO |  |

Program, constants, calls, libfs, etc.

| STNO | DC | 437 |
| :---: | :--- | :--- |
| PRNO | DC | 4 | | Your student number |
| :--- |
| Problem number |
| Program, constants, |
| Calls, libfs, etc. |

Before each machine language instruction is executed, AM tests the instruction to determine if its execution would alter the core resident monitor, alter AM itself, make an invalid entry to a subroutine, or an invalid alteration of a subroutine. If its execution would have one of these undesireable effects, further execution of the student's program is suppressed and a link is made to DBUG as explanned elsewhere. Further execution of the program is also suppressed if a valid exit is reached, or the run time becomes excessive.

If, on the other hand, AM decides that the instruction should be allowed to execute, the instruction counter is incremented and control is passed to the instruction that was tested.

Immediately after the execution of the instruction, control returns to AM by means of a hardware interrupt. This interrupt results from the machine being in interrupt run (also called trace) mode. AM then tests the next instruction, as before. This procedure of first testing each instruction and then allowing its execution is continued until further execution of the program is suppressed, as described above.

To cause the Assembler Monitor System to monitor your program, you need only call AM at the beginning of your program. When control is passed to AM, it reads the student's input data from the disk, initializes parameters to be used during execution to tell how core has been partitioned for the core load, forces the operator to place the machine in interrupt run mode, and gives control to the testing portion of $A M$ so as to test the first instruction of the student's program.

## Interpretation of Output

After the Assembler Monitor has decided that the student's program should not be allowed to execute further, control is passed to DBUG. DBUG moves the paper to the top of the page and prints on the right-hand side the student number, problem number, contents of the accumulator, extension, index, carry, and overflow registers and the floating accumulator. Student number and problem number are given as positive decimal numbers; the accumulator, extension, and index registers are given in hexadecimal; the carry and overflow are given as being "on" or "off"; and the floating accumulator is given in hexadecimal and decimal.

On the left-hand side is printed a core map which gives the starting addresses and lengths of eleven consecutive partitions that make up a core load. The lengths of these partitions vary according to the program(s) in the core load.

The first partition is the Index Register Area, which consists of the first four words of core (i.e., addresses 0 , 1, 2, and 3). It is so called because it includes the three index registers, which are in words 1, 2, and 3 in core.

The second partition is the resident monitor, which includes the core resident monitor supplied by IBM (excluding the first four words of core) and the core image header which is located immediately thereafter.

The third partition is the mainline program, which includes everything from the end of the core image header to the beginning of the Assembler Monitor (AM).

The fourth partition is the AMS program, which consists of the program AM, and is the in-core part of the Assembler Monitor System.

The fifth partition is the subroutine area, which includes all subroutines, regardless of type, located between the AMS program and the interrupt level subroutine area.

The sixth partition is the interrupt level subroutine area, which includes all interrupt level subroutines except levels two and four.

The seventh partition is unused core. This partition of core is not used by the core load.

The eighth partition is the LIBF transfer vector, which consists of three words for each library function entry point in the core load.

The ninth partition is the floating accumulator, which consists of six words of core used as an accumulator for floating point arithmetic. There is no floating accumulator if there is no LIBF transfer vector.

The tenth partition of core is the CALL transfer vector, which consists of one word for each CALL entry point in the core load. The CALL transfer vector will sometimes include a dummy word in order to make the floating accumulator begin on an even core boundary.

The eleventh and last partition of core is COMMON, which is located at the very end of core. It is in this partition of core that the input and output occur. COMMON is saved between LINKs by the monitor system; i.e., it is still in core when DBUG and GROUT are loaded in turn.

On the left-hand edge the starting address and length of each partition are printed in hexadecimal. On the righthand edge the word ADDR is printed beside that partition in which the effective address of the instruction causing the exit was located. If the exit was not caused by the effective address, the word PREA is printed beside the partition in which the last effective address formed was located.

DBUG then skips a space and prints the instruction causing the exit and the prior instruction in hexadecimal. To the left it prints the real address (the address of the instruction in core) and the loading address (the address of the instruction relative to the loading point of the mainline, which is the address found on a relocatable assembler mainline listing or a FORTRAN mainline listing).

If the program failed to clear location \$IOCT (/0032 hexadecimal), a line is printed indicating this fact. This error would indicate that an interrupt service subroutine was not incrementing or decrementing \$IOCT properly. Location \$IOCT should be zero if and only if there are no $1 / 0$ interrupts pending.

A line is then printed giving the reason why the student's program was prevented from further execution, i.e., the reason for exiting. This line is printed in the form:

AMS $x x$ (message giving reason for exit) where $x x$ is the error number. The error numbers are given in the following table:

```
0 0 \text { Instruction is located in COMMON.}
01 Instruction is located in CALL transfer vector.
02 Instruction is located in floating accumulator.
03 Instruction is located in LIBF transfer vector illegally.
0 4 \text { Instruction is located in unused core.}
0 5 \text { Instruction is located in interrupt level subroutine.}
0 6 ~ I n s t r u c t i o n ~ i s ~ l o c a t e d ~ i n ~ s u b r o u t i n e ~ a r e a ~ i l l e g a l l y .
07 Instruction is located in AMS program.
0 9 ~ I n s t r u c t i o n ~ i s ~ l o c a t e d ~ i n ~ m o n i t o r ~ i l l e g a l l y . ~
OA Instruction is located in index register area.
OC Attempt to alter CALL transfer vector.
OE Attempt to alter LIBF transfer vector.
10 Attempt to alter interrupt level subroutine.
ll Attempt to alter subroutine area from mainline.
12 Attempt to alter AMS program.
14 Attempt to alter resident monitor.
15 Attempt to alter word zero in core.
1A 64 instructions did irrelevent access of core.
1B Program terminated due to excessive run time.
1C Invalid instruction.
20 Valid exit.
```

Any other indicators indicate an error in the Assembler Monitor System, and should not occur.

Next, the message ADDRESSES OF LAST n. INSTRUCTIONS EXECUTED is printed, where $\underline{n}$ is a decimal number with a maximum value of 64 giving the number of addresses listed thereafter. If the program ran for less than or equal to 64 program steps; all the addresses, in the order of execution, will be listed. If the program ran for more than 64 program steps, only the addresses of the last 64 are listed. Both the real and loading addresses are listed in hexadecimal.

If any instructions did an irrelevent access of core (i.e., they did no harm, but did no good, either), then the addresses of these instructions come out in a table titled ADDRESSES OF INSTRUCTIONS LOADING IRRELEVENT DATA where
is a hexadecimal number. As above, each address is given both relative to the beginning of core ("REAL") and relative to the beginning of the core load ("LOAD").

Ir the event that the problem number is zero (or is not the number of a defined problem) no LINK is made to GROUT, the program is not graded, and the only other information printed is the program load length (both in hex and decimal) and the number of instructions executed (both in hex and decimal).

If the problem number is that of a defined problem, then a link is made to GROUT which outputs the student's grade and reasons behind it in three sections titled POINTS FOR CORRECT ANSWERS, ADDITIONAL POINTS FOR OUTPUT, and POINTS FOR PROGRAM EFFICIENCY. The total grade is the product of the total points for each of the three sections (divided by one million to scale it down). The total points for each section is printed after the word total at the bottom of each section and is equal to the sum of the points earned under that section as listed under the right-hand column. The points for each line are calculated from how well the student program did relative to the standard on this point. The total grade is printed beside the message TOTAL GRADE EOUALS at the bottom of the page. The total grade and each of the separate totals should range from zero to one thousand, although it is not impossible to make a grade greater than one thousand.

After printing the total grade, control is returned to the IBM supplied monitor supervisor, which begins looking for the next job.

Operator Procedure and Interpretation of Operator Console Displays

With student program decks in the card reader and the system initialized, the console typewriter will display the following message:

SET MODE SW TO INT RUN
At this time, the operator must set the mode switch (located on the right hand side of the display panel) to "interrupt run" and press the "program start" switch. If "program start" is pressed without first: setting the "interrupt run" condition, the above message will be printed again. If the machine is already in the interrupt run mode, the message will not be printed. While in interrupt run mode, the "stop" button will have no effect.

The Assembler Monitor System has a provision for terminating a student due to excessive run time (based on a count of operations executed) and this is done automatically. However, an operator may abort a student program by momentarily placing the bit ll switch on the console in the up position. In case this does not abort the program and cause an appropriate error message to be printed, then the program is not under Assembler Monitor System control.

If an abort is desired while the machine is in the interrupt run mode and not under control of the Assembler Monitor System, the operator must first take the machine out of interrupt rin mode and then press "interrupt request." Alternately, he can first press "interrupt request" which will stop the computer, then change to the run mode and press "program start."

If bit switch 0 is up, the program will stop after each machine language instruction is executed under control of the Assembler Monitor System and display the contents of the Accumulator, Extension and Carry and Overflow status.

Bit switches 14 and 15 are used to control student core dumps and displays to the operator during student program execution under control of the Assembler Monitor System. If bit switch 14 is up and 15 is down, all relevant student core content will be dumped on the printer and the system will pass to the next student Erogram. If bit switch 15 is up, the computer will pause and display a coded error number in the storage buffer register, the
address of the instruction causing the exit in the accumulator, and the effective address of the last instruction employing an effective address in the extension register. Upon restarting, if bit switch 14 is also up, then the relevant student core data will be dumped on the printer. With neither switch 14 or 15 up, no pause or dump occurs.

An override feature is provided which may be used with caution: if bit switch 13 is up after a pause caused by a program exit and switch 15 being up then the Assembler Monitor System will return to the student program.

Programs, Subroutines and Files
Runining of the student programs is done under the control of the Assembler Monitor System. This system consists of seven main computer programs, several standard subroutines and four data files described briefly below.

The Assembler Monitor Program (AM) serves as a direct monitor over the running of the student's program, with each instruction performed under monitor control. A debugging aid generator program (DBUGT) prints out a trace and other diagnostic aids to the student from information provided by the AM. The raw grade is calculated by a grading program (GROUT) which calculates the students grade, prints it and records it for the instructor. program GRINT generates information on which the grade is based from the standard problem supplied. Program INITD initializes data for the grading of each student's problem. For the start of a grading run or for each new problem set, the system is reinitialized with program RINIT which clears the data and grade files. A message input program (MSGIN) loads file a message file with the appropriate messages to be used by the DBUGT program.

Subroutines used in the system include the IBMsupplied Commercial Subroutine Package-Version III, and assembler subroutine for floating binary to decimal (FBTD) and the following special subroutines: FORMT and SHIFT are used by DBUGT to decipher assembler instructions HEXIN converts four alphanumeric characters representing a hexadecimal number into the integer equivalent. HEX and HXOUT convert an integer back to hexadecimal. DCOUT converts an integer into five alpha characters representing a number in decimal. OUT prints a line of alphanumeric characters and clears the output buffer to blanks. DSCTR dumps a 320 word core sector (length of one disk sector) in hexadecimal to the printer.

SAVGR contains three records of 320 words per record. Since each disk sector contains 320 words, this file uses three sectors. The initial contents of SAVGR are unimportant because AM loads the file with new data with each new student program. The accual instructions, variables, and constants of AM are stored by that program in three blocks. The three records are the 320 words following respectively the three DSA statements labeled IOARI, IOAR2, and IOAR3. It is the task of program DBUGT to extract the pertinent data from irrelevant coding. SAVGR is referred to in all programs by symbolic file number 1.

MSGBF also contains three records of 320 words per record, giving three disk sectors. It is used by program DBUGT to print all words interpreting the output of $A M$ including all headings and in converting all numbers from integer format to alpha characters. To initialize MSGBF, program MSGIN is executed, reading data from twelve cards in FORMAT (80Al), and storing the contents on disk. Refer to program listing of MSGIN for contents of data cards. MSGBF is referred to in all programs by symbolic file number 2.

The records of GFILE each contain 16 words with one record generated per student program run under the system for grading. The length of GFILE can therefore be varied with the needs of the user by simply changing the number of sectors specified when the file is set up and by changing the number of records in the DEFINE FILE statement in program GROUT. For example, if the user desired 400 records at 20 words per sector, this would require 20 sectors of disk. The contents of each record of GFILE will be listed and explained later. The contents of GFILE is initially set to zeroes by program RINIT. GFILE is referred to by symbolic file number 3.

DATFT contains ten records of 320 words apiece, giving 10 disk sectors. Each record contains information used by the system in grading a problem of the standard data set. The system can therefore handle a problem set of 10 problems. The corresponding record of DATFT must be reset to zeroes before entering a new standard problem in the problem set. To reset. DATFT and/or GFILE, execute program RINIT, following it by one data card of FORMAT (10I2,10x,I2). The first 10 fields indicate which records of DATFT are to be reinitialized. If GFILE is also to be 32 is to be left blank. DATFT will be referred to by symbolic file number 4.

To define these four files on disk, the computer should be given instructions corresponding to these:
// JOB
// DUP
*STOREDATA WS FX SAVGR0003
*STOREDATA WS FX MSGBF0003
*STOREDATA WS FX GFILEOO20
*STOREDATA WS FX DATFTOOlO

Since programs DBUGT, GRINP, and GROUT are executed by links and have quite lengthy core-loads, the running of a student program under the Assembler Monitor System can be quite time-consuming. If the user has sufficient area on disk, it is suggested that these programs be stored CoreImage. This will considerably speed the operation of the system. All four data files must therefore be stored in Fixed Area on disk.

The next step is to execute program MSGIN which will read 12 cards of alphanumeric data and initialize file MSGBF (see program listing). This file will be used to generate headers and output information by program DBUGT.

Assembler Monitor Use
The Assembler Monitor System has provision for up to 120 words of input data read by the student program determining the grade on up to 120 words of output. Tine input is loaded by AM into COMMON, beginning with the last word of core. AM will not load input data beyond the end of a student's specified COMMON. Any COMMON beyond the number of words of input is filled with zero or some other easily recognizable "garbage word" specified by the instructor. This is done as a debugging aid so that the student can determine by examining a core dump what, if anything, his program has changed. The output must also be in COMMON and within the last 320 words of core. The 120 words of output can be divided into as many as 10 blocks of consedutive core locations and these blocks can be located anywhere within COMMON. This permits freedom to:

1. Give more important answers more credit for grade.
2. Court part of the grade on intermediate answers arrived at in the process of generating the final answers.
3. Remove points for destroying the input in the process of obtaining an answer. A further option is provided to give points for partially correct answers, that is answers either in the correct blocks but in incorrect order, or answers found anywhere within COMMON. This option can be used as a debugging aid by pointing out to the student that he has made only a small logic error in addressing and not written a program that does nothing.

Program efficiency is determined on the basis of five parameters: mainline program length, subroutine length, length of COMMON, number of instructions executed, and a standard curve or bias. The curve is based on the theory that with the high speed of this computer, the length of most programs run under the system, the difficulty of writing in assembler, and inexperience in proframing of most students using the system, that a program that works should not receive a failing grade no matter how inefficient it is.

In order to initialize DATFT with the standard input data, output buffer locations, and grading factors the instructor must perform the following operations: First Store subroutines INITD, HEX and HEXTN on disk. HEXIN is
used to translate core addresses entered in hexadecimal (four characters) into integer constants. HEX is used to translate DATFT to hexadecimal characters for dump to printer. INITD takes parameters problem number and standard input and data cards for output locations and grading points and puts them on disk. Since INITD is a subroutine, it cannot initialize its own HO. This must be done by a short calling program(written in FORTRANj. This program must initialize ISS routines for disk, card reader, and line printer and must tell. INITD where to find DATFT on disk. For Example;

## //ЈøB

$/ /$ F $\varnothing$ R

* $\varnothing$ NE WøRD INTEGERS
*EXTENDED PRECISIØN
*IøCS (LISK,CARD, 1403 PRINTER)


DEFINE FILE $4(20,320, \mathrm{U}, \mathrm{K})$
-program. (see below) --

CALI INITD (... ) CALL EXIT END
//XEQ 01
*FILES ( 4, DATFT)
-
-
5 Data cards.
The following four integer calling arguments should be passed to INITD if called by FøRTRAN:

1. Problem Number (PROBN).
2. Standard Input (STDIP), the first element of an array up to 120 words long.
3. Stanc ard Input Length (STDIL), the number of words of input.
4. "Garbage" Word (GBGWD), filler for remaining student COMMON; e.g.. CALL INITD (PROBN,STDIP, STDIL, GBGWD) •

The array STDIP can be initilized by data statements, arithmetis assignment statements, or read statements in integer or Al format. (Do not use the commercial CALL READ.) If it is desired to place real numbers into STDIP it must be remembered that one extended precision real number fills three words of core and that the first element of a real array should be equivalenced to the third element of the corresponding integer array. This is because FORTRAN arrays are stored in reverse order in core. For the same reason, the first element of STDIP will be placed by AM into the last word in core, and following elements will be stored into aescending core locations.

If greater versatility of input is desired, the FøRTRAN program can call an assembler subroutine which generates STDIP and in turn calls INITD. In this way, the student can be provided with input in the format of actual instructions, characters in card-code, paper tape, etc. These changes in the calling sequence must be noted: All calling arguments must be addresses of the parameters, not the parameters themselves. Also, STDIP is the address of the last location of input. For example:

|  | ENT |  | DATA |
| :---: | :---: | :---: | :---: |
| DATA | - . |  | -•• |
|  | CALL |  | INITD |
|  | DC |  | PROBN |
|  | DC |  | STDIP |
|  | DC |  | STDIL |
|  | DC |  | GBGWD |
|  | EXIT |  |  |
| PROBN | DC |  | 1 |
| STDIP | BES | E | 120 |
| STDIL | DC |  | 120 |
| GBGWD | DC |  | /EEEE |
|  | END |  |  |

In this way, STDIP can be filled by such assembler pseudo-ops as:

| DC | hex constant |
| :--- | :--- |
| DEC | 2-word decimal integer or real constant |
| XFLC | entended precision real constant |
| EBC | entended BCD interchange code characters |
| DMES | printer hex (console, ll32 or l403) |
| DN | name code constant. |


#### Abstract

The instructor can provide, by an LIBF to ZIPCO, paper tape or card-code characters.

Output locations and grading parameters are entered as data on five cards after the //XEQ and *FILES cards (and also after any data cards read by the mainline). The first two cards contain respectively the beginning and ending addresses of up to ten output buffers GROUT is to search for answers. The addresses are to be expressed in four digit hexadecimal, absolute, with two spaces between address, up to ter addresses per card.

FØRMAT (10(4A1,2X)) Card three contains five numbers which are the percentage points to be assigned for program efficiency. The first number is for mainline program length, the second for subroutine length, the third for length of COMMON, the fourth for number of instructions executed, and the fifth is the curve. The sum of all five parameters should equal 100. Each number should be expresses as three digits with two spaces between each.

FøRMAT (5 (13,2X)) Card four contains up to ten percentage points for answers in correct locations, one corresponding to each answer buffer defined in cards one and two. Card five contains three percentage points determining value of partially correct answers. The first parameder is percentage for completely correct answers, the second for answers within the correct buffers but not necessarily in correct order, and the third is for answers anywhere within COMMON. The sum of cards Gour and five must each equal 100. The formats are the same as for card three. for example (for a machine with 8 K . core):


| IFFO | IFDQ | IFEO lFE8 |
| :--- | :--- | :--- | :--- |
| IFFF | IFDF | IFE7 $1 F E 8$ |


| 010 | 010 | 010 | 040 | 030 |
| :--- | :--- | :--- | :--- | :--- |
| 015 | 015 | 040 | 030 |  |
| 070 | 020 | 010 |  |  |

At the end of execution, INITD will give a hex dump of DATFT to the printer. The standard input buffer is stored in DATFT in reverse order to that in which it is loaded into core. The first element of DATFT (last element in the $F \not \subset R T R A N$ dump) is loaded into the last location of core and so forth.

## Initilization of Standard Programs

The final step in preparing the system for grading student programs is to run the standard programs. These are to be run in the same manner as student programs, with the following changes in operating procedure:

1. Parameters to be passed to AM are the address of problem number and a student number of -1 (FFFF in hexadecimal).
2. All data switches on the console must be placed in the up position (FFFF hexadecimal).
3. The program will stop after the first instruction with an exit code of 301C hexadecimal in the SBR. All switches except 13 should be placed in the down position and the program started. The program should now stop with 3020 in the SBR (normal Exit). (If a core-dump is desired, put switch 14 up.) Restart the machine. AM will now store the information it has compiled on SAVGR to DBUGT, which will read SAVGR, MSGBF, and DATFT. DBUGT will determine that the program is a standard and will link to GRINP. GRINP will complete the initilization of DATFT with standard output and standard program efficiency. COMMON and DATFT will be dumped to the printer in hexadecimal. A link will be performed back to DBUGT, which will then handle the standard as if it were a normal student program (as a cross-check on the standard.) The standard program will receive a grade of 1000 points. All student programs will be graded in comparison to this standard grade. Student programs can now be run and graded on the system for all problems on which the standard has been initialized.

## Computation of Grade

The computation of the student's grade is based on these factors:
I. Answers
A. completely correct
B. partially correct
II. Program efficiency
A. Mainline length
B. Subroutine length
C. Length of COMMON
D. Number of instructions executed
E. Standard curve
III. Correct termination of program (EXIT)

To compute $I, A$, GROUT compares the contents of the output data blocks in the students COMMON to the corresponding standard output block, and computes the ratio of the number
of correct answers the student finds to the length of the block (standard number of correct answers). This ratio is multiplied by the corresponding grading parameter for correct answers (entered into DATFT by INITD, data card \#4). The sum of these 10 products is then multiplied by the grading factor for totally correct answers (INITD, first number, data card \#5). GROUT then searches the student's output buffers, counting the number of correct answers placed anywhere within the correct data block. The ratio of the number of answers so found to the total number of possible answers, is multiplied by the grading parameter for answers within the correct data blocks (INITD, second number, card \#5). All of COMMON is then searched for the correct answers found in any locations, the ratio to total answers is computed and multiplied by the parameter for answers within COMMON. The total points for answers is the sum of points for correct answers, answers within the correct buffers, and answers anywhere in COMMON.

Points for program efficiency are computed as the sum of points for program length, subroutine length, length of COMMON, number of instructions executed and standard curve. Points for program length are computed as the ratio of Standard program length to student program length, times the grading parameter for program length (INITD, first number, card \#3). If the student did not receive a perfect score on answers and his program length was less than that of the standard, points for program length is computed as if his program length was the same as that of the standard. Points for subroutine, COMMON, and number of instructions are computed in a like manner.

Total grade is computed by multiplying points for answers by points for program efficiency. $25 \%$ of the grade is lost if the program is terminated by anything but a standard exit (AMS 20). A message to this effect is printed. The final grade is then scaled on a factor of 1000. It is important to note that the grade given by the system is based upon a comparison between the student program and a "standard" program, and not between the student and other student programs. For this reason, the final scaling of grades must be left to the instructor. The system does, however, give a fair grade in that the grade is proportional to the worth of the program (if the grading parameters are assigned properly), and that the instructor can easily tell from the output supplied to him, where to scale the grades.

## Output of GROUT to GFILE

GROUT supplies certain pertinent information about the student's grade to the instructor by entering a 16 word record on GFILE for each program graded, unless the student passes a negative student number to AM. The

```
contents of GFILE is as follows:
    1. Record number (first record has total number of
        records saved).
    2. Student number.
    3. Problem number.
    4. AMS exit code.
    5. Total grade.
    6. Points for completely correct answers.
    7. Points for all answers.
    8. Points for program efficiency.
    9. Program length.
10. Subroutine length.
11. Length of COMMON.
12. and 13. Number of instructions executed. Since
        a program can possibly execute more than 32,767
        instructions, (the greatest possible i:iteger the
        machine can hold), AM divides the instruction
        count into two words. The first is the number
        of instructions divided by 10000, and the second
        is the remainder of the instruction count. In
        other words, }13\mathrm{ is the low order four decimal
        digits and }12\mathrm{ is the.upper decimal digits.
    14. Number of answers in correct locations.
    15. Number of answers within correct data blocks.
    16. Number of answers anywhere within COMMON.
```


## CONCLUSION AND RECOMMENDATIONS

Difficulty with a fixed problem set to be usea repeatedly, led to the approach employed which permits new problems to be introduced as frequently as necessary. This has been effective over several quarters. Experience has shown that a first program for the student should be extremely simple - something like reading a number into the computer and printing it out. This divorces the mechanics of basic input and output from other programming complexities and gives the student the satisfaction of having been on the computer very early in the course.

Additional instructions have been considered for the repertoire of the simulator. These might include arithmetic and cyclic shifts, multiplication and perhaps even division. Although these would permit the solution of more sophisticated problems and may make the simulated computer more like an actual one, they would not make a major advance to the learning obtained via the current basic machine commands.

Provision is made in the present systems for accommodating the five decimal digit student identification number at Florida Institute of Technology. This is inadequate for some schools and will ultimately be inadequate at F.I.T. when a change to Social Security numbers as identification occurs, as it most surely will.

The Assembly Monitor system is only serving a small quantity of people - those computer science majors who use it in machine language programming. However, they are not required to use it. Moreover, nearly all problems at the machine language level, have been individually designed and must result in a working program. Further work on this program is not recommended at this time.

10 T AAA LOAD ACCUMULATOR - LDA
The contents of the Accumulator are replaced by the contents of the effective address. The contents of the effective address are not changed. The Sign latch is set equal to the sign of the contents of effective address. The Overflow latch is not affected.

$$
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A \text { ) }
$$

Example: $104625 \quad$ EA=625+213=838


Example: ll 0001 EA=001

Before execution:
Accumulator -999999
Location 001 ???????
Sign Latch

After execution:
Accumulator -999999
Location 001 -999999
Sign Latch Negative

The contents of the specified Index Register $T$ are replaced by the contents of the effective address AAA. The contents of the effective address are not affected. The Sign latch is set equal to the sign of the contents of the effective address. The Overflow latch is not affected.
$E A=A A A$ (Note: $T$ cannot be 0 ; this instruction must specify an Index Register.)

Example: $409123 \quad \mathrm{EA}=123$
Before execution: After execution:

I/R 9 ???????
Location 123 -999995
Sign Latch

After execution:
I/R 9 -999995
Location 123 -999995
Sign Latch Negative

20 T AAA ADD TO ACCUMULATOR - ADD
The contents of the effective address are algebraically added to the contents of the Accumulator. The contents of the effective address are not changed. The sign latch is set equal to the sign of the result in the Accumulator. The Overflow latch is set on if sum exceeds +999999 or is less than -999999. When Overflow occurs, high-order digits are truncated. The Overflow latch is setOFF if overflow did not occur.

$$
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A \text { ) }
$$

Examples:
Accumulator Before EA -999999 -000001 $-001001+000001$ $-999999+999999$ $+010010000000$ $+999999+000001$
$+999999 \quad+999999$


Accumulator After 000000
-001000
000000
$+010010$
+999998

Over-

| flow |  | Sign |
| :--- | :---: | :---: |
| ON |  | 0 |
| OFF | Neg. |  |
| OFF | 0 |  |
| OFF | + |  |
| ON | 0 |  |
| ON | + |  |



The contents of the effective address are algebraically subtracted from the contents of the Accumulator. The contents of the effective address are not changed. The sign latch is set equal to the sign of the result in the Accumulator. The Overflow latch is set on if the result is greater than +999999 or less than -999999 . When overflow occurs, high-order digits are truncated. The Overflow latch is set off if overflow did not occur.

$$
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A \text { ) }
$$

Examples:
Accumulator Before EA
-999999 -999999
-999999 +000001
$+000001+999000$
+999998 +000001
+999999 -999999


## 42 T AAA ADD TO INDEX REGISTER - MD

The contents of the effective address AAA are algebraically added to the contents of the specified Index Register $T$. The contents of the effective address are not changed. The Sign latch is set equal to the sign of the result in the Index Register. The Overflow latch is set on if sum exceeds +999999 or is less than -999999. When overflow occurs, high-order digits are truncated. The Overflow latch is set off if overflow did not occur.

```
EA = AAA (Note: T cannot be 0; this instruction must
``` specify an Index Register.)

Example: 425002

Before execution:
\begin{tabular}{lr} 
Before execution: \\
\hline I/R 5 & -999999 \\
Location 002 & -000001 \\
Sign latch & \(?\) \\
Overflow latch & \(?\)
\end{tabular}

After execution:
\begin{tabular}{ll} 
IR/5 & 000000 \\
Location 002 & -000001 \\
Sign latch & 0 \\
Overflow latch ON
\end{tabular}

\section*{41 T AAA STORE INDEX REGISTER - STX}

The contents of the effective address AAA are replaced by the contents of the specified Index Register \(T\). The contents of Index Register \(T\) are not affected. The Sign latch is set equal to the sign of the contents of Index Register T. The Overflow latch is not affected.
\(E A=A A A\) (Note: T cannot be 0; this instruction must specify an Index Register.)

Example: \(411402 \quad E A=402\)
Before execution:
I/R \(1 \quad 000000\)
Location 402 ??????
Sign Latch

60 T AAA READ A CARD - IN
Data is read in from a card and temporarily held in a buffer area. The data in the buffer is then checked for validity. If the first column contains an asterisk, the current program is terminated. If not, the first column must be a blank, plus sign, or minus sign. Blank is treated as a plus sign. Columns 2 through 7 must contain digits from 0 to 9 --- blanks are not allowed. Columns 8 - 80 may contain comments.

If the validity checking does not detect an error, the data is loaded into the core location specified by the effective address. If the data is invalid, the contents of the effective address are not altered. The Overflow and Sign latches are not affected in any case.
\[
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A)
\]

Example: 601427 (data in card, +426351) \(E A=427+111=538\)

Before execution:
I/R 1 +000111
Location 538 ??????

After execution:
I/R 1 +000111
Location 538 +426351

The contents of the effective address is printed on the printer, and the paper is advanced one space. The Sign and Overflow latches are not affected.
```

77 0 000 STOP - HLT

```

Execution is terminated. The Sign and Overflow latches are not affected. Core is dumped onto the printer, ten locations per line for any line containing a word in which any change has been made in storage during execution.

50 T AAA BRANCH (Unconditional) - B
Control is transferred to the instruction at the effective address. The Sign and Overflow latches are not affected.
\[
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A)
\]

Example:
\begin{tabular}{cl} 
Core location & Contents \\
\hline 042 & 500862 \\
043 & \(3 ? ? ? ? ?\) \\
962 & 210044
\end{tabular}

Execution of the Branch instruction at location 042 will cause the next instruction executed to be the subtract instruction at location 862 .

\section*{51 T NZA BRANCH NEGATIVE - BN}

This instruction causes a branch to the effective address if the Sign latch is Negative. If the Sign latch is not negative, control goes to the next sequential address. The sign and Overflow latches are not altered.
\(E \dot{A}=A A A+\) contents of \(T\) (if \(T=0, E A=A A A\) )

\section*{52 T AAA BRANCH ZERO - BZ}

This instruction causes a branch to the effective address if the gign latch is zero. Otherwise, control. goes to the next sequential address. The Sign and Overflow latches are not altered.
\[
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A)
\]

\section*{53 T AAA BRANCH POSITIVE - BP}

This instruction causes a branch to the effective address if the Sign latch is positive. Otherwise, control goes to the next sequential address. The Sign and Overflow latches are not altered.
\[
\mathrm{EA}=\mathrm{AAA}+\text { contents of } \mathrm{T} \text { (if } \mathrm{T}=0, \mathrm{EA}=\mathrm{AAA})
\]

54 T AAA BRANCH OVERFLOW - B \(\varnothing\)
This instruction causes a branch to the effective address if the Overflow latch is ON. Otherwise, cortrol goes to the next sequential address. If branch occurs, then the Overflow latch is reset to OFF. The Sign latch is not affected.
\[
E A=A A A+\text { contents of } T \text { (if } T=0, E A=A A A)
\]

\section*{APPENDIX II}

Problems l, 3, 4, 5, and 6 are from the winter quarter 1969. Problems \(11,12,13,14\) are from the spring quarter 1969.

PROBLEM NC. 1
Given: A set of 100 data cards containing values \(X_{i}\) such that:
\[
\begin{aligned}
\mathrm{i}= & 1,2,3, \ldots, 100 \\
& -999999 \leq \mathrm{X}_{\mathrm{i}} \leq+999999
\end{aligned}
\]

Write a machine ianguage problem beginning in location 0 (zero) to solve the following equation:

100
\(\sum X_{i}\) where \(0 \leq X_{i} \leq 1000\)
\(i=1\)
i.e.; omit values of \(X_{i}\) outside of the above range from the sum.

Be as efficient as possible.
Write out the answer on the printer. Store your answer in location 900. Read the given input data into locations 500-599.
Use index register(s) and conditional instruction(s).

\section*{PROBLEM NO. 3}

Given: A set of 100 data cards containing values \(X_{i}\) such that:
\[
\begin{aligned}
& i=1, \quad 2,3, \ldots 100 \\
& -999999 \leq X_{i} \leq+999999
\end{aligned}
\]

Write a machi 2 language program beginning in location 0 (zero) to solve the following equations:
\[
\begin{aligned}
& \text { Sum } 1=\sum_{i=1}^{99} X_{i} \quad \begin{array}{c}
\text { (Sum the contents of only the odd } \\
\text { numbered locations: } i=1,3,5, \ldots .99)
\end{array} \\
& \text { Sum } 2=\sum_{i=2}^{100} X_{i} \quad \begin{array}{l}
\text { (Sum the contents of only the even } \\
\text { numbered locations: } i=2,4,6, \ldots 100)
\end{array}
\end{aligned}
\]

Write out both answers on the printer.
Store the answers: Sum 1 in location 900
Sum 2 in location 901
Read the given input data into locations 500-599. Assume no overflow will occur.
Use any instructions you think necessary. Be as efficient as possible.

PROBLEM NO. 4
Given: A set of 100 data cards containing values \(X_{i}\) such that:
\[
\begin{aligned}
& i=1,2, \quad 3, \ldots 100 \\
& -999999 \leq X_{i} \leq+999999
\end{aligned}
\]

Write a machine language program beginning in location 0 (zero) to perform the following:
(a) Find ANS. \(1=\) total number of negative items in the list
(b) Find ANS. \(2=\) total number of zero items in the list.
(c) Find ANS. \(3=\) total number of positive items in the list.

Problem No. 4 (cont'd)
Read the given data into locations 500-599
Store the answers: Ans. 1 in loc 900
Ans. 2 in loc 901
Ans. 3 in loc 903
Use any instructions you think necessary. Be as efficient as possible.

PROBLEM NO. 5
Given: Two sets of 50 data cards containing values
\[
\begin{aligned}
& \left.\begin{array}{l}
X_{i} \\
Y_{i}
\end{array}\right\} \quad i=1,2,3, \ldots 50 \\
& \text { such that } \begin{array}{l}
-5000 \leq X_{i} \leq+5000 \\
\\
-5000 \leq Y_{i} \leq+5000
\end{array}
\end{aligned}
\]

Find the sum of the differences \(\left(X_{i}-Y_{i}\right)\) by the following formula:
\[
\sum_{i=1}^{50}\left(X_{i}-Y_{i}\right)
\]

Read the first set of fifty cards into locations 500-549.
Read the second set of fifty cards into locations 550-599.
Write out the answer on the printer.
Store the answer in location 900.
Use any instructions you think necessary. Be as efficient as possible.

PROBLEM NO, 6
Determine and print the first \(N\) numbers of the "FIBBONACCI" series. In the "FIBBONACCI" series each number is the sum of the previous two numbers with the first two numbers of the series being 0 and 1 .

Example of the "FIBBONACCI" series:
\(0,1,1,2,3,5,8 \ldots(\) to \(N\) terms of the series)
Read the value of \(N\) into location 500. Store the terms of the series starting in location 900. Print the terms of the series.

PROBLEM NO. 11
Write a program which will evaluate
\[
f(x)=3 x^{2}+2 x+7
\]
for \(x\) an integer \((0<x<100)\) to be read in from a data card. Test \(x\) after reading to make sure it is correct. Print out the value of \(x\) and \(f(x)\). Store \(f(x)\) in 900. If the value of \(x\) is out of the allowable range, print out the actual value of \(x, 000000\) for \(f(x)\), and stop.

PROBLEM NO. 12
Write a program which will read (1) a card with the integer \(0<N<100\). (2) \(N\) data cards into \(N\) successive locations, then sort the \(N\) numbers into ascending order and print them out. Read the data cards into locations \(200 f f\) and sort into locations \(300 f f\).

PROBLEM NO. 13
Write a program to read in 25 numbers. These are to be stored in consecutive locations starting at 200. The numbers represent consecutive elements in consecutive rows of a matrix. Perform the transpose of the matrix so that rows and columns are interchanged. Print out the trans posed matrix. Store transpose in locations \(300 f f\).

PROBLEM NO. 14
Given three sets of data cards of \(N \leq 30\) cards each: Read the first set of \(N\) cards into locations 100, 103, 106,...

Read the second set of \(N\) cards into locations 101, 104, 107,...

Read the third set of \(N\) cards into locations 102, 105, 108, ...

Print out in order locations 100, 102, etc. \(N\) is on first card. (A total of \(3 N+1\) cards will be read.)

\section*{APPENDIX III}

This appendix contains summaries of the results of three surveys conducted after the automated problem sets had been used by several classes.

First is the student response to a questionaire which followed the course.

Second is the concensus of the instructor who taught the course.

Third is observations of the IBM-1130 operator who actually accepted the students programs and batch processed them.

\section*{Student Survey on Automated Problcin Sets}

A questionnaire (Table I) was prepared to ascertain the effectivity of the automated problem sets from the standpoint of the students. This questionnaire and the summarized responses from 134 students are shown. The questions were designed to determine the extent of ease or difficulty which the new (to the students) concept of machine language was assimilated. Results were obtained after the student had subsequently been exposed to, and had written programs in, a compiler language, namely, FORTRAN.

The final question requesting comments on improvement of the course elicited response from approximately fifty percent of the questionnaires. It opened a Pandora's box with a great diversity of opinions expressed. At the extremes, these ranged from the ideas that machine language was a complete waste of time and all programming training should be concentrated on FORTRAN to the desire to have the full quarter devoted to binary machine language with more emphasis on arithmetic and control unit organization. Specific comments also dealt with insufficient demonstration on keypunch, need to have first programs examined in detail by instructor before attempting to run, need for monitors to be better versed in the simulation language and in the problems assigned that quarter. A majority of the opinions expressed reflected the students' personal desires in results of such a course and in their success or frustrations in achieving these desires.

The following numbered observations correspond to the questions of the same number shown in Table \(I\).
1. Less than two percent of the students had any prior experience with machine language.
2. Eighty percent believed that the instruction set was about the right complexity with the rest equally divided between too simple and too complex.
3. Responses were equally divided between those accepting the set as adequate and those desiring a multiply instruction. The fact that a negligible number thought shifting should be included probably indicates that its use was not pointed out to the students.
4. A negligible number of responses felt that the number of branch instructions was excessive and about a third wanted even more variations.
5. Opinion was about 7-5 in favor of a less restrictive I/0 set.
6. Opinion was about equally divided for and against inelusion of logical instructions.
7. The decimal coding was almost universally accepted as suitable for grasping the essentials of machine langage. A few dissidents identified a desire for binary.
8. Less than twenty percent considered the brief study of machine language a waste of time for the ultimate user.
9. All debugging aids provided proved helpful but the greatest aid was discussion with other students.
10. Difficulties with getting ultimately successful runs were most impeded by the actual closed shop mechanism of the Computer Center (probably underqualified monitors, bugs still in the program, and general lack of understanding of procedures). Failure to understand the function of the simulated computer operations and errors in card punching were also substantaal contributors.
ll. A large majority (over ninety-eight percent) considered the problem set reasonably difficult with the rest equally divided between too hard and too easy.
12. 12 . Problem difficulty was rated roughly equal.
13.) The most difficult problems took three quarters of the
14. \(\}\) students less than four hours of homework and less than five computer runs.
16.\} The easiest problem took three quarters of the students
17.) less than two hours of homework and less than three computer runs.
18.) Results of this question appear to belie the preceding
19. two results. For if the program were indeed tested and ready for the run for record it should succeed on the first, or at worst, second run. The statistics indicate that many used four or more of these runs on their more difficult problem.
20. A majority felt that there was a sufficient diversity in the problem set although several felt the problems were too similar.
21.\} Analysis, coding and debugging difficulty varied much 22. between individuals and no one stood out as uniformly particularly hard or particularly easy.


TABLE I
TO: Students who took CSI62 during Winter Telm 1969
FROM: D. R. Clutterham, Head of Mathematical Sciences Dept.
We need to obtain some information regarding the use of the simulated computer used to teach machine language in the CSi62 course. Please complete the following questionnaire as accurately as possible and return to the Mathematical Sciences Department in person or by campus mail. If desired you may delete the portion above the double line to preserve anonymity. Please complete and return immodiately.

Underline answer which fits your case.
1. Had you ever worked with machine language before?
(a) yes (b) no
2. The instruction set provided was
(a) too complex,
(b) about right",
(c) too elementary
3. The arithmetic instructions
(a) were adequate, (b) should have included shifting,
(c) should have included multiplication.
4. The branch instruction set
(a) was adequate, (b) could be improved with some additional types, (c) had too many alternatives.
5. The input/output set of instructions was
(a) too restrictive
(b) adequate
(c) should permit formatting
6. Logic instructions should be included
(a) no, (b) such as "AND", "OR", "COMPLEMENT."
7. Greater understanding of machine language would have been obtained if numbers and codes had been
(a) in octal, (b) in hexadecimal, (c) in binary,
(d) the decimal used was adequate.
8. The study of machine language
(a) is a waste of time for an ultimate user
(b) gave me a much better understanding of computers
(c) contributed to my appreciation of FORTRAN
9. The most helpful debugging aid was
(a) the program trace, (b) the memory and status dump,
(c) discussion with monitor, (d) discussion with
classmates
10. The greatest difficulty in completing a program successfully was
(a) incomplete understarding of instructions
(b) getting results from a run on the comprter
(c) punching an accurate set of cards
17. The problem set to be solved
(a) was adequate, (b) was too difficult,
(c) was too easy
12. The problem which was most difficult for me was
(a) 1
(b) 2
(c) 3
(d) 4
(e) 5
13. The problem which was most difficult for me required
(a) less than 2 hours of homework
(b) two to 4 hours of homework
(c) four to li hours of homework
(d) over 10 hours of homework
14. The problem which was most difficult for me required
(a) less than 3 computer runs, (b) 3 to 5 computer runs
(c) 6 to 9 computer runs, (d) more than 9 computer runs
15. The problem which was easiest for me was
(a) 1
(b) 2
(c) 3
(d) \(4 \quad\) (e) 5
(f) don't remember
16. The problem which was easiest Eor me required
(a) less than 2 hours of homework, (b) 2 to 4 hours of homework, (c) 4 to 10 hours of homework
(d) over 10 hours of homework
17. The problem which was easiest for me required
(a) less than 3 computer runs, (b) 3 or 4 computer runs
(c) 5 to 7 computer runs, (d) 8 or more computer runs
18. My easiest problem ran correctly on my run for record
number
(a) 1
(b) 2
(c) 3
(d) 4 or greater
19. My hardest problem ran correctly on my run for record number
(a) 1
(b) 2
(c) 3
(d) 4 or greater
20. The problems in our problem set
(a) were about right, (b) were too similar,
(c) were too different
21. The part of these problems I found easiest was (a) analysis, (b) coding, (c) debugging
22. The part of these problems I found hardest was (a) analysis, (b) coding, (c) debugging
23. Include any comments for improving this part of the course.

Suxvey of Instructors Using the Automated Problem Sets
Seven instructors have been introduced to the automated problem sets and five have taught the introductory computer course at Florida Institute of Terhnology using the sets. Their observarions are sunimarized here.

When a class is given a common problem, there is a tendency to either copy the solution of one of the better students or to work collectively on a program so that the net result is several groups of identical solutions. This problem is not pecuiiar to this course or even this subject, but usually students vary their own solutions from the one they copy and this is not. done with the automated problem sets. One solution may be to have the students turn in their handwritten coding sheet before they begin their actual machine debugging; then their final programs should be modifications to the handwritien ones. Another solution is to develop a very large set of similar problems so that students have essentially an individual problem.

The instruction code set seems generally suitable to the instructors. More experienced instructors found the set quite suitable or else desired only a shifft operation. Newly indoctrinated instructors desired a multiply and perhaps also a divide instruction. Somewhat more capability in the input-output format appears desirable, although exactly what form it should take was not agreed upon. A set of left and right shifts with and without e circular capability have been designed for the program but are not now incorporated.

One anticipated problem - that of teaching the use of the keypunch in classroom - has not arisen; learning the use of the keypunci seems to be passed very readily between the students, and a minimum of words from the instructor is sufficient.

Survey of Machine Operators Using the Simulator
An initial complaint of the operators was that instructors did not sufficiently define the problem to the students and further definition had to be supplied in detail. This is recognized as a continuing problem and the instructors are putting more care and detail into the definition.

A second difficulty is that assignments are relatively few, but everyone's problem comes due at the same time. Even if the assignments are given well in advance, normal student procrastination causes a heavy run on both the card punching equipment and on the computer in the last couple of days before a grading run is due. A solution to this problem, as yet untried, is to stager problem due dates giving easier problems to the students whose problems are due first. In addition to this, simpler problems could be given much earlier in the quarter so that the students can first learn some of the mechanics of preparing a problem ior the machine and getting basic input-output mastered.

APPENDIX IV



```

// JOB LOAOPOOL
// * LOAOPOO2
// * PROGRAM TO LOAD EACH STUDENT PROGRAN INTO PSEUDO CORE, AND LOAOPOO3
// * BRING IN THE FILE DF STANDARD DATA ON THE PRCBLEM FOR GRADING. LOADPOO4
/1 *
// DUP
L,0900 LOAOPOO6
// FOR
*OELETE LOAOP LOAOPOO7
*NAME LOADP
*LIST SDURCE PROGRAM
*LIST SUBPROGRAM NAMES
*LIST SYMBOL TABLE
*IOCS(CARO,OISK,1403 PRINTER)
*EXTENOEO PRECISION
*ONE WORD INTEGERS
INTEGER ERRS
INTEGER ERROR
INTEGER LOC(2000),XR(28),AREG(2),TAG,ADDR,EA,OPCOC,NEUNC(2)
INTEGER IOBUF(4B),NAME(32),ERRCT(5)
INTEGER TABLE(16)
INTEGER RNTIM(2),PRDGL
INTEGER LOCl(1000),LOC2(1000),XR1(9),XR2(9) LOAOPO22
INTEGER NSAVI(30),NSAV2(30)
INTEGER STORT,STOPL
INTEGER ANS1(30),ANS2(30),NANS,LCANS(5),NANSR(5)
INTEGER NROSR(lO),LOCRO(lO)
INTEGER PTSR,PTSA,PTSW,PTS
INTEGER FOATA,POSPT(3)
INTEGER PTCR(10),PTCRN,PTCA(10),PTCC(10),PTCO,PTCW(10),PTHO
INTEGER PCGRT,PCGPL
INTEGER ROATA(14)
INTEGER FILND,PC,LINE(70),OATA(212),DATA1(106),OATA2(106)

```

```

    COMMON LOC,XR,AREG,ISIGN,INSTR,TAG,ADOR,EA,DFCOC,NEUMO,TOBUF,NAME
    COMMON ERRCT
COMMON NI,NO,TABLE,JERR
COMMON I,J,K,L,M LOAOPO36
COMMON INIT
COMMON NSTUD,NPRDB
COMMON RNTIM,PROGL,NDCDS
COMMUN NANSW,NSAVI,NSAV2
COMMON IDUMY,STDRT,STDPL,ANS1,ANS2,NOROS,NRGPS,POSPT
GOMMON IDUMY,STDRT,STDPL,ANS1,ANS2,NOROS,NRGPS,POSPT
COMMON PTHO,NANS,FDATA,MAXRT,PCGRT,PCGPL,ROATA LOAOPO43
COMMON PTSR,PTSA,PTSW,PTS
COMMON FILNO,PC,IOVFL,LINE,OATA
EQUIVAL.ENCE (LOC(1),LOCI(1)),(LOC(1001),LOC2(1)) LOAOPO46
EQUIVALENCE (XR(1),XR1(1)),(XR(10),XR211))
EQUIVALENEE (POSPT(1),NPPTR),(POSPT(2),NPPTAI,(POSPT(3),NPPTK)
EQUIVALENCE (NI,NOCOS), (NWTR\&NANSW)
EQUIVALENCE SOATAl(l),OATA(l)),(DATAZ(.l),OATA(107))
EQUIVALENCE (LOCll,LOCl(l)O,(LOCl2,LOC2(l))
EQUIVALENCE (ERRCTII) ERROR)
DEFINE FILE l(24, 160,U,NXREC)
C------TEST FOR MONITOR CARO
1 IF(ERROR) 2,10,10
2 IF(IOBUF(2)-TABLE(1)) 10,20,10
10 EA=1
CALL RCR60
GO TO l
C------SKIP. TO NEW PAGE, PRINT MONITOR CARO
20 FRITE(NO,22) IDBUF
22 FORMAT(lHL,16Al,32A2,//)
00 27 I=2,15
OO 26 J=1,10
LOADPOO4
LOAOP007
LOAOPOO8
Name Loadp
LOAOPOO9
LOAOPO10
LOAOPOL2
*IOCS(CARO,OISK,1403 PRINTER) LOAOPOI3
LOAOPO14
LOAOPO15
LOAOPO16
LOAOPO18
LDADPO19
LOADPO20
LOAOP022
LOOPO26
LOAOP027
L0A0P02\&
LOAOPO29
LOAOPOSO
LOAOP031
LOAOP032
LOAOP033
COMMON L,NO,TABLE,JERR LOAOPO35
LOAOPO36
LOAOP037
LDADP038
LOADPO39
LDAOP042
LOAOPO43
LOAOP044
LOAOP045
LOAOP046
LOADP047
LOADPO48
LOADP049
LOAOPO51
LOAOP051
LOAOPOS3
LOAOP054
LOADP055
LOADP055
LOAOPO56
LOADPO57
LOAOPO58
LOAOPO59
LOADP059
LDAOP060
LDAOPOG1
LOADP062
LOAOP063
LOAOPO64

```
```

        IF(IOBUF(I)-TABLE(J)) 26,27,26 LOACPO65
        26 CONTINUE LOAOP066
    C---ERROR IOBUF(I) SET TO ZERD. LOAOP067
J=1
27 IOBUF(I) = J-I
C------TEST FOR MONITOR START CARD
NCARD = (()(IDOUF(2)*10+IDBUF(3))*10+IOBUF(4))*10+IDBUF(5))*10+
1IOBUF(6))*10+10BUF(7)
LDADPO72
NCARD = ((IOBUF(4)*10+IOBUF(5))*10+IOBUF(6))*10+IOBUF(7) LOADPO73
NO LIST OF SOURCE PROGRAM IF NCARD EQUALS ZERO. LOADPO74
IF(NCARD-1) 28,29,10 LOALPO75
28 NCARD = 2
LOAOP076
29 NOCOS = 1 LOADPC77
NSTUD = (()IOBUF(9)*10+1DBUF(10))*10+10BUF(11))*10*10BUF(12))*10+1LOADPO78
lOBUF(13) LOACPO79
NPROB = 10*IOBUF(14) + IOBUF(15) LOADPO8O
DO 45 I= l,32 LOAOPO81
45 NAME(I) = {DBUF(I+16) LOADPO82
c

```



```

C MONITOR CARDS ARE IDENTIFIED BY AN ASTERISK IN COLUMN 1. LOADPO87
C ROUTINE RETURNS ON READING A MONITOR CARD, OR WHEN CORE LDAC LOACPOB
C EXCEEOS PSUEOD-CORE.
C IOBUF CONTAINS LAST RECORD READ. LOADPOSI
C ERRS CONTAINS COUNT OF ERROR FLAGS. LOADPOG2
C ERRORS ARE FLAGGED WITH AN ASTERISK ON LISTING. LOACPOG3
C ERRGRS ARE ALHAYS LISTED.
C LOADING STARTS IN CDRE LDCATION ZERO. LOADPOG5
c clear psuedo-core.
LOREDODO97
AREG(1)=2.0000 LOADPO98
AREG(2)=20000 LOADPO99
OO 3 IAR=1,18 LOAOPIOO
3 XR(IAR) = 25000
O0 106 IAR=1,2000 LOACP102
106 LOC(IAR)=30000
INITIALIZE IAR AND ERRS. LOADPIO4
IAR=0
ERRS=0
110 EA = IAR + 1
CALL RDR6O
IF(ERROR) 120,140,130
130 ERRS=ERRS+1
PUT ASTERISK IN ERROR FLAG.
ERROR=TABLE(15)
GO TO }15
BLANK DUT ERRDR FLAG
140 ERROR=TABLE(12)
NO LIST DF SOURCE PROGRAM IF NCARD EQUALS ThO
GO TO(150,160),NCARD
150 WRITE(NO,51) ERROR,IAR,IOBUF
51 FORMAT(1H,AL,IX,I4,4X,7Al,4X,9A1,32A2) LOADP119
160 IAR=1AR+1 LDADP120
C TEST FOR ENO OF PSUEDD-CORE.
IF(IAR-999) 110,110,120
C-----* CARO OR ENO OF CORE ENCOUNTERED LOADP122
LOADP123
120 WRITE(NO,51) TABLEIl2I,IAR,IOBUF LCACP124
PROGL=IAR LOACP125
ERROR = 0
C-----ABORT IF MISPUNCHED CARD IN DECK.
IF(ERRS) 30,30,10
LOADP127
30 IF(IOBUF(3)-TABLE(1)) 31.2.31
LCACP129
LCADP130

```






```

    16C1 K=NCCCS + FLATA - 1 SIMRN329
    LOCI(EA) = DATAI(K) SIMRN330
    LGCZ(EA) = CATA2(K) SIMRN331
    GO TC 5CO SIMRN332
    1605 CALL RDR60 SIMRN333
    IF(EKRCT(1)) 385,50C,385 SIMRN334
    385 JERR=4
    GO TC 5CO SIMRN336
    C
C
1610 CALL CECEB(CEAR,KKBLF)
WRITEINC,1615IKKBUF
1615 FORMAT(1H ,7A1)
NWTR=NWTR+1
IF.(NhTR-30) 1627,1617,500
NSAVI(NWTR)=CEAR1 SIMNSM
NSAV2(NHTR)=CEAR2
GO TC 500
C

```

```

177C IAR = PC
JERR=1
GO TO 500 O..: SIMRN354
C SIMRN35S
C----------------------------------------------------------------------------- SIMRN357
c
SIMRN35
C
C
SIMRN359
TCNTR=TCNTR+1
1F(TCNTR-25) 510,510,501
501 CALL CATSW(1,J)
GO TO (510,520),J
510 NSW=1
C
IF(LCTR) 570,560,570
560 LCTR=1
MRITUNC,561) SIMRN369
WRITE(NC,561) SIMRN37G

```


```

C-----GET C(ADOR) SIGN (ANFL \
570 CALL CECEB(MREG,IIBUF)
SIMRN374
SIMRN375
NEUMC(1)= NUTBL(2*INSTR-1) SIMRN376
NEUMO(2)=NUTBL(2*INSTR) SIMMN377
C----GET C(XR) SIMRN378
IF(TAG) 580,580,585
580 [F(IAG) 580,580,585
JJBLF(I)= TABLE(12) SIMRN381
MMBLYF(I) = TABLE(12) SIMRN382
522 CONTINUE
GO TO 590
C----GETC(XR) SIMRN378
SIMRN379
SIMRN38O
52 CONTINUE
58 CALL DECEB(CXR,JJBUF)
585 NXRGI = XRI(TAG) SUF)
NXRG2 = XR2(TAG)
CALL DECEB(NXREG,MMBUF)
C-----GET C(EA)
590 CALL CECEB(CEAR,KKBUF)
NNREG(1)=LOC1(EA)
NNREG(2)=LOC2(EA)
CALL CECEB(NNREG,NNBUF)
C-----GET C(ACC)
SIMRN382
SIMRN383
SIMRN384
SIMRN386
SIMRN387
SIMRN389
SIMRN390
SIMRN391
SIMRN392
SIMRN392
SIMRN394

```
```

CALL DECEB(AREG,LLBUF)
SIMRN395
C
EA=EA-1
WRITE(NC,596) RNTIM(2),PC,IIBUF,NEUMD,TAG,ACDR,JJBUF,EA,KKBUF
LLBUF,MMBUF, NNBUF, ISIGN,IOVFL
SIMRN396
SIMRN397
SIMRN398
SIMRN399
596 FDRMAT(1H , 15, 2X,14,4X,7A1,4X,2A2,I1,1X,13,4X,7A1,4X,14,3X,7A1,5X,SIMRN400
C 7A1,4X,7A1,5X,7A1,6X,I2,5X,12 ) SIMRN401
C SIMRN4O2
GO TO 521
SIMRN4O3
C
C
SIMRN404
520 GO TO(512,521),MSW
SIMRN407
512 MSW=2
WRITE(NO,555)
55 FDRMAT(1H)
SIMRN408
521 IF(JERR) BOO,523,800
521 IF(JERR) BOO,523,800
SIMRN411
C----FLUSH NC NEXT JOB (SIMGIO) IF SSW 11 ON SIMRN412
C I OPERATOR JUDGES TIME EXCESSIVE -IF PRINTING IN LODP SIMRN4I3
WILL NOT BE STCPPED IN REASONABLE TIME BY CCUNTER.
SIMRN414
523 CALL DATSW(11,J)
SIMRN415
GO TO (530,600),J
C
SIMRN416
SIMRN417
C-----BEGIN NEXT MACHINE CYCLE
SIMRN41B
600 PC=IAR SIMRN419
C----FLUSH TC NEXT PROGRAM (AFTER DUMP) IF RUN TIME EXCESSIVE. SIMRN42O
IF(RNTIM(2)-MAXRT) 1000,1000,530 SIMRN42I
530 JERR=3 SIMRN422
800 RETURN SIMRN423
ENO
// DUP SIMRN
SIMRN424
SIMRN425
*DELETE SIMRN
SIMRN426
*STORE WS UA SIMRN
SIMRN427

```


```

C---N-NANS = NO CF ANSWERS REGUIREC ULNPG131
C DATA GARD G, hORD 4. GUMPGI32
724 DO 7265 Kl=1,5 CLNPG133
LI=1
7241 IF(Ll-NANSR(K1)) 7242,7242,7265
7242 IAR = LCANS(Kl) + LI DLMPGI36
LI=L1 + 1 . ULMPG137
IF(LOCI(IAR)-ANS1(I)) 726,725,726
725 1F(LOC2(IAR)-ANS2(1)) 726,727,726
726 GO TG 7241
7265 CONTINUE
GO TC 728 DUNPGL42
OLNPG134
7241 IF(LI-NANSR(K1)) 7242,7242,7265 DLMPG135
DLMPG13\&
DUMPG139
DLMPG14C
DLMPGL41
C------PTCO = NO OF POINTS FOR CORRECT ANS IN CORR LCCS IN ANY CRCER.DLMPGI43
CATA CARD 9, WGRD 2 OLMPGI44
727 PTSA= PTSA+ PTCO DLMPG145
7 2 8 GO TO 7101 DUMPG146
730 CONTINUE
C
C------PTS = TOTAL NG. CF PQINTS RECIEVEC.
PTS = PTSR + PTSA + PTSW
C
NPPTT = NO. OF POSSIBLE PCINTS - TCTAL.
NPPTT = NPPTR + NPPTA + NPPTW
PPTT = NPPTT
RWGTM = L.
RWGPL = 1%
IF(PTS-NPPTT) 742,750,750
C------DO NCT COUNT TIME OR LENGTH BETTER THAN STANOARL IF FULL
POINTS WERE NOT EARNED.
742 IF(RNTIM(2)-STCRT) 744,744,743
743 RWGTM = RWGTM*STDRT/RNTIM(2)
744 IF(PROGL-STCPL) 760,760,745
745 RWGPL = RWGPL*STDPL/PROGL
GO TO 760
750 RWGTM = RWGTM*STDRT/RNTIM(2)
RWGPL = RHGPL*STDPL/PROGL
760 RAWGR=(100-PCGRT-PCGPL+PCGRT*RWGTM+PCGPL*RWGPL)/PPTT*10.*PTS
GO TO(780,770,770,780),JERR
770 RAWGR = 3*RAWGR/4
780 CONTINUE
C
C-----ROUTINE TO CUNP PSUEDO-CORE TC PRINTER.
LOC IS 1000 WORC PSUEDO-CDRE.
DUMP IS TEN 7II INTEGERS PER LINE.
ALL OF CORE IS DUMPED.
WRITE(NO,799) NAME,NPROB
799 FORMAT(1HO,08X,32A2,12X,'PRCBLEN NC.',I4)
IF(NPROB-4)7995,8205,7995
E205 NANS=0
7995 GO TG(801,803,805,807),JERR
801 WRITE(NO,802)
802 FORMAT(1HO,'EXECUTICN COMPLETE')
GO TO 820
803 WRITE(NC,804) PC
804 FORMATIIHO,'EXECUTION TERMINATEC BY INVALID INSTRUCTIGN AT ',I3)
GO TO 820
805 WRITE(NC,806)
806 FORMAT(IHO,'EXECUTION TERNINATED DUE TO EXCESSIVE RUN TIME') DUMPGI92
GO TC 82O
8 0 7 WRITE(NC,808) PC,EA DUMPGL94
DUMPG193
808 FORMAT(IHO,'EXECUTION TERMINATEC BY INSTR. AT ',I3,' ATTEMPTING TODUMPG195
I REAC IST CARD OF NEXT PROG. INTO I,I3)
DUMPG196

```
```

    \becausei乚 NRIIEINC,G11) RNTIN(2),STLRT,PRCGL,STOPL,NCGCS,NCRCS,NANSW,NANS DUMPG19I
    &11 FORNAT(1HO,IOX,'RLNTIME',14X,'LENGTH CF CECK',C8X,'NO CF CARES , DUMPGIG8
    l,'iLEAC',OGX,'NC OF ANSWERS hRITTEN'/4(05X,'YCURS',C6X,'STANCARC'),DUMPGlg9
    2/,3x,8(16,C6x)/1
    5 FORNAT(IHO,C3X,'POINTS RECEIVED FOR---I/O5X,'REACING DATA',llX,
    1 'ANS IN CORK LCCATICNS',C4X,
    <'WRITING ANSWERS',C9X,'TCTAL',19X:'RAW',/
    ```

```

    34(OSX,'YOURS',C6X,'STANDARC'),4X,'GRADE',/,3X,9(IE,6X)/) DUMPG205
    CALL CECES(AREG;KEUFF)
    WRITE(NC,813) ISIGN,ICVFL,K8UFF
    ```

```

C゙-----PRINT INDEX REGISTERS.
I AR = C
J=8
CO 860 K=1,9
IAR =IAR + I
REGI=XRI(IAR)
REG2=XR2(IAR)
C-----CLEAR UNUSED INDEX REGISTERS
IF(REGI-25000) 831,832,831
\&32 IF(REG2-250:`0) 831,833,831
833 CO 834 L=1,7
LINE(J)=TA8LE(12)
834 J= J+l
GO TC }86
831 CALL DECE8(REG,LINE(J))
J=J + 7
860 CONTINUE
GRITE(NO,843) (LINE(J),J=8,70)
843 FORMAT(1H ,!0X,5HI/RS ,9(3X,7A1),/%
C-~--DUMP PSEUDO - CORE.
I AR = 0
DO 8jC I=1,10C
J=1
M=0
OD B\&C K=1,10
IAR =IAR + I
REGl=LOCl(IAR)
REG2=LOC2(IAR)
IF(REGI-300CO) 851,852,851
852 IF(REG2-300CO) 851,853,851
853 DO 854 L=1,7
LINE(J)=TA8LE(12)
J=J+l
M=M+1
GO TG 880
B5l CALL CECEB(REG,LINE(J))
J=J + 7
C
CO NCT PRINT LINE IF ALL LCCATIONS IN IT UNAFFECTED BY PROGRAM.
CO NCT PRINT LINE I
821 J=IAR-1C
WRITE(NC,822) J,LINE
822 FORMAT(1H,I3,2X,10(3X,7A1))
830 CONTINUE
C----EERRCR TRAP
C
C-----IF FINAL GRAUE RUN, WRITE GRADE INFC ON FILE.
C-----IF INITIALIZATIDN,GCTC INILG, IF STUD. PRDG. GD TD LOAD NEXT. DUMPG257
IF(INIT) 885,881,890. DUMPG258
881 READ(2'l) NFILE
NFILE = NFILE + 1 OUMPG259
DUMPG259
W'RITE(2'1) NFILE DUMPG261
WRITE(2'NFILE) NPRDB,NSTUC,JERR,RNYIM(2),PRCGL,PTSR,PTSA,PYSW.NAM.DUMPG262

```
```

                                    RAWGR
                                    DUMPG263
    8 8 5 \text { CALL LINK(LCADP) DUMPG264}
    890 CALL LINK(INI2G)
            END
                    DLMPG266
                    DUMPG267
    \#DELETE DUMPG DUMPG268
\#STORECI WS UA DUMPG 0001
*FILES(2,SMSTU)
DLMPG260
DUMPG270

```
```

1/ \# INITGCOI
//*PROGRAM TC INITIALILE GRADER. INITGOG2
l/*PROGRAM TC INITIALILE GRADER.
// FOR
*NAME INITG
\#IDCS(CARD,DISK,1403 PRINTER)
\#EXTENDED PRECISIDN
*ONE WORD INTEGERS
*LIST SDURCE PROGRAM
*LIST SUBPROGRAM NAMES
*LIST SYMBDL TABLE
INTEGER A(2205),INPUT(160),CROIN(78),NREM(77),CATA(212)
INTEGER NRDSR(1C),TABLE(16)
INTEGER NROSR(1C),TABLE(16)
INTEGER DATAl(106),DATA2(106)
INTEGER FOATA
COMMON A,INPUT, NREM,OATA
EQUIVALENCE (NPROB,A(2140)),(CRCIN(1),INPUT (69)),(TABLE(1),A(2116)INITGO1B
1),(INIT,A(2138)),(NROSR(1),CROIN(1)),(NCRES +INPUT(64)),(NRGPS,INPUINITGO19
l);(INIT,A(2138)),(NROSR(1),CROIN(L)),(NCRES+INPUT(64)),(NRGPS,INPUINITGOI9
EGUIVALENCE (EA,A(2C25)) INITGO2I
EQUIVALENCE(LCCI2,A(1)),ILOC12,A(1001)) INITG022
EQUIVALENCE (OATAI(1),DATA(1)),(OATA2(1),CATA(lO7)) INITGO23
EGUIVALENCE {ERR,A(2109))
EQUIVALENCE (NI,A(2114)),(NC,A(2115))
EQUIVALENCE(FOAIA,CROIN(75))
DEFINE FILE 1(24,160,U,NXREC)
CEFINE FILE S(12,106;U,NXRDC)
1 INIT = 1
CO \& 1=1,160
8 INPUT(I)=0
NI=2
NO=5
REAU(NI,13) TABLE,NDTST
13 FORMAT(16A1;I1) INITGO35
CALL OATSW(3,J) INITGO36
GO TO\5CO,1C),J
500 REAC(NI,II) NPROB,NRCSR(I),FDATA INITGO38
GD TG 600
IO REAC(NI,L1) NPROB INITGO4O
C REAO(NI,LI) NROSK,LECRO,LCANS,PTCR,PTGA,PTCC,PTCH,PTCRN,PTCC, INITGC4I
C LPThC,NANS,FDATA
REAC(NI,ll) CRCIN
600 CONTINUE
11 FORMAT(10({6,2X)1
CO 2CI = 1,10
NRGPS=I-1
K = MROSR(I)
IF(K) 2C,21,2C
20 NORCS = NORDS +K
NRGPS = 10

```

```

ILOCRL,LCANS,PTGR,PTCA,PTCC,PTCW,PTCRN,PTCC,PTWC,NANS,FOATAOPCSPT
21 HRITE(I'NPROB) INPUT
21 HRITE(1'NPROB) INPUT
14 IF(NUTST-6) 15,15,16 INITGG56
15 REAL{512*NDTST-1) CATA
GO IC 19
16 EA=1
16 EA=I I I 1.106
GALL KCRUC
CAT\DeltaI(I)=LOC11
l:ATAZ(I)=LQCl2
IF(ENR) 3,2,3
TGCO3
INITGCO4
INITGOOG
INITG007
INITGOOB

* IST SOURCE program
INITG009
NNITGOI
INITGOLO
INITC...ili
NNITGCI2
INITGC12
INITGOI3
NITGO14
INITGO15
INITGOl6
6)INITGO1B
INITG024
INITG025
INITGO26
DEFINE FILE l(24,160,U,NXREC) INITG027
NITGO27
INITGC29
I=1,160 INITGO30
INITGO3I
INITGC32
INITG033
INITGO34
GO TO\5CO,101,J
INITGO38
INITGO39
INITGO40
INITG042
INITGO42
INITGO43
INITG044
NTG044
*)
INITGO45
INITGO46
INITG047
NITG048
INITG048
INITGO49
INITGOSO
INITGOS2
INITGC53
INITGO54
INITGO54
INITGCS6
. INITGC57
. INITGG58
INITGOS9
INITGC6C
INITGC6C
INITGCGI
INITGC62
NITGC62
(ENR) 3,2,3 = INCI2 INITGC63

```
```

        3 PAUSE 7C09 INITG065
        I = I - I
        LOC12 = IABS(LCCl2)
        4 WRITE(NC,17) LOC11,LOC12
    17 FORNAT(1H,14,13)
    19 CALL RDR60
        CALL LINK(LOADP)
        END
                                    INITG066
    INITG067
NITG068
INI TG069
INITG070
END
// DUP INITG
// DUP INITG
*STORECI hS UA INITG 0001
*FILES(l,FSTDG),(5,SIMDT)
INITG071
INITGC72
INITG073
NITGO74
INITG075
INITGO76

```

```

    730 CONIINLE
                                    IN12GCE5
        NANS = NANSn INI2GC66
        IF (NANS) 77C,770,740 INI2GC67
    740 CO 76C I=1,NANS - INJ2GCEE
    ANSI(I) = NSAVI(I) INI2GC69
    ANS2(I) = NSAV2(I) INI2GC7C
    IF(I-30) 760,77C,770 INI2GC71
    7EO CONTINUE
    770 STORT = RNTIN(2)
        STDPL = PRDGL
        CALL LINKIDUMPGI
    C------P:IT RESULTS OF SECOND PASS THRU CUMPG INTC STC VARIABLES
C AND PUT STANDARD DATA ON F!LE.
101 NPPTR = PTSR
NPPTA = PTSA
NPPTW = PTSK
F WRITE(ITNPROB) IDUMY,STGRT,STDPL,ANSI,ANS2,NCRCY:\RGPS,NRDSR,
C LLOCRO,LCANS,PTCR,PTCA,PTCC,PTCW,PTCRN,PTCC,PTWC, AANS,FDATA,PCSPT
HRITEIIINPROBI INPUT
PAUSE 3333
C RETURN TO INITIALIZE ANOTHER PRCBLEN IF SENSE SWITCH 2 ON.
CALL DATSW(2,J)
GO TO(2CO,777),J
200 CALL LINK(INITG)
777 STOP 7777
END
// DUP
*DELETE INI2G
\#STORECI WS UA INI2G OOOL
*FILES(l,FSTDG)
INI2GC95

```
```

// JCB . ROSTDOOl
// FOR
RDSTDOO2
*LIST ALL
*ONE hCRD INTEGERS
\#EXTENCED PRECISICN
RCSTDC03
ROSTDCO4
RDSTOCOS
SUBRCUTINE RESTC
INTEGER A(22C5),INPUT(16C),NREM(75)
COMNCN A,INPUT,NREN
EQUIVALENCE(NPRCB,A(214C))
EQUIVALENCE (NAXRT,INPUT(144))
\imath
C-----THIS RCLTINE READS THE fILE MADE FRCM THE STANDARC FOR THE
C FROELEM THE STUCENT IS ATTEMPTING. RCSTDOI3
MAXRT = 500C
IF(NPRCB) 5,5,2
2 IF(NPRCB-24) 10,1C,5 ROSTDC17
5 DO \& = l , 1C3 ROSTDOL8
INPLT(I) = C
MPRCB = 0 C ROSTDC2O
RETURN
10 READ(I'NPROB) INPUT ROSTDC22
RETLRN ROSTDC23
END
// DUP
*DELETE RCSTC RDSTDO26
*STQRE hS UA RCSTC
RCSTDOO6
RESTDC07
RDSTDOO\&
RESTDCO9
RDSTDOlO
RCSTOCll
RCSTDC11
C
RCSTOO13
RCSTDC14
l
RDSTDO15
RDSTDC16
RDSTDC19
RDSTDC21
ROSTDC22
RDSTDC23
RCSTDC24
RDSTDO25
RDSTDO26
RDSTDC27

```
```

/1 JOB % RCR6OCCL
// DUP RCR60CO2
// FOR
RCR60CO3
RRCR60C04

* LIST ALL
*EXTENCED PRECISICN
RCR60CO5
RDR6COC6
*ONE WORD INTEGERS
SUBRDUTINE RORGC
INTEGER KBUFF(7)
KBUFF(7) RRR60C08
INTEGER LOC(2C00),XR(18),AREG(2),TAG,ADCR,EA,DPCCC,NEUNC(2) RDRGOCO9
INTEGER IOBUF(80), ERRCT(5)}\quad\mathrm{ RCRGOCIC
INTEGER TABLE(16)
INTEGER TABLE(16) RDRGOCII
INTEGER LDC1(1COO),LOC2(1000),XR1(9),XR2(9) RERGOC12
INTEGER ERRCR RRRGCCIJ
COMNDN LOC,XR,AREG,ISIGN,INSTR,TAG,ACCR,EA,CPCCC,NEUMC,IOBUF,ERRCTRCR6OC14
COMNDN NI,NC,TABLE - RCRGOCL5
EQUIVALENCE (LCC(1),LOC1(1)),{LEC(1OC1),LCC2{1)) RERGOC16
EQUIVALENCE (XR(1),XR1(1)),(XR(10),XR2(1))
EQUIVALENCE (XR(1),XR1(1)),(XR(10),XR2(1)) R RCR60C17
C
C
C ROUTINE TO SIMULATE REAC INSTRUCTIEN
ON RETURN--ERKCR IS SET -I IF ASTERISK CARC REAC RORGOC22
ON RETURN--ERRCR IS SET - I IF ASTERISK CARC REAC
C IF NO ERRCR
RCRG0C }2
RCR60C24
PSUEDO-CORE LDCATIDN IS NOT ALTERED IF ASTERISK CARC IS REAC, IF R RCRGOC25
RER60C26
RRCRGOC27
RCR60028
RER60C29
ERRCR=0
READ(NI,II) (IDBUF(I),I=1,48) R RCR60C3C
ll FORMAT(16A1;32A2)
NOCDS=NOCDS+1
RETURN IF MDNITCR(ASTERISK) CARC.
IFIIOBUF(1)-TABLE(15)) 30,20,30
20 ERRCR=(-1) . RRR60C35
RETURN
C CONVERT 7AI TD 7II.
C CONVERT 7AI TD 7IL.
30 DO 21 N=12,14
IF(TABLE(N)-IGBUF(1)) 21,25,21
.l CONTINUE
IF(TABLE(11)-IDBUF(1)) 22,23,22
\angleL ERRGR=1
GO TD 50
23 KBUFF(1)=-1
GD TO 26
25 KBUFF(1)=1
26 DD 29 N=2,7
OD 2% J=1,10
MO 28 J=1,10
28 CONTINUE
GO TO 22
29 KBUFF(N)= J-1
C PACK 7II INTO 2I3 AND STORE INTC PSUEDO-CCRE.
LOC1(EA) = ({KBUFF(2)*10+KBUFF(3))*10+KEUFF(4))*KRUFF(1) ROR60C55
LOCZ(EA) = ({KBUFF(5)*10+KBUFF(6))*10+KEUFF(7))*KEUFF(1))
5 0 ~ R E T L R N
ENO
RCR60057
RCR60C58
// DUP
*DELETE ROR60
\#OELETE WS UA ROR60 RDR60 % R R OR60C6C
MEAD(NI,11) (IDBUF(I),I=1,48)
RCR6CC 32
C

```

```

    OR IF INVALID OATA IS READ.
    C
C
C
C
C
RCR60C07
RERGOC19
RCR60C20
RDR60033
RCR60034
RDR60C35
RDR60036
RCR60C 37
RCR60C38
RCR60039
F(TABLE(N)-IGBUF(1))21,25,21 RCR60C4C
M- RCR60042

```

```

    RCR60C45
    RCR60C46
    RCR60C47
RCR60C48
RCR60C49
RDR60050
28 CONTINUE . N
RCR60051

```

```

    RCR60C52
    C
RDR60053
RER60C54
RDR60059
M RCR60C6C
RCR60C6C
RDRGOG61

```
```

M,

```
DECEBCI2
CECEBCl3
DECEBCl4
DECEBC14
DECEBO15
DECEBC16
DECEBO17
DECEBO18
DECEBO 19
CECEBC20
DECEBC21
DECEBO22
DECEBG23
DECEBC24
DECEBC25
DECEBC26
DECEBO27
CECEBC28
DECEBC29
decebo3o
DECEBO31
DECEBO32
DECEBO33
DECEBO34

```

// JOB 0015
// FOR LATCHOOL
*LIST ALL
*EXTENCED PRECISIDN
\#ONE GORD INTEGERS
SUBROUTINE LATCH(REG)
INTEGER REG(2)
INTEGER COREI2C2O:
COMMON CORE,ISIGN

```

```

C THIS ROUTINE SETS THE SIGN INOICATOR, ISIGN, TC -I,O,+1
ACCORDING TO THE SIGN DF THE DATA IN A.
EXAMPLES...
REG(1) REG(2) ISIGN
-999 000 - -1 DATA IS NEGATIVE
000 -999 -1 DATA IS NEGATIVE
000 000 0 CATA IS ZERO
CATA IS ZERC
DATA IS POSITIVE
IF ( REG(1) ) 30,20,50
20 IF ( REG(2)) 30,40,50
30 [SIGN=-1
RETURN
40 ISIGN=0
RETURN
50. ISIGN=1
RETURN
END
// DUP
*DELETE LATCH
*STORE WS UA LATCH

```

LATCHOOL
LATCHCO2
LATCHCO 3
LATCHCO4
LATCHOOS
LATCHOO6
LATCHCOT
LATCHCO8
LATCHOO9
LATCHOIO
LATCHOLI
LATCHOI2
LATCHOL 3
LATCHO14
LATCHOL 5
LATCHCl6
LATCHCI7
LATCHCl8
LATCHOL9
LATCHO2O
LATCHCZ1
LATCHC22
LATCHC23
LATCHC 24
LATCHC25
LATCH026
LATCHC27
LATCHC28
LATCHC 29
LATCHC 30
LATCHO31
LATCHC 32
LATCHC 33
LATCHO34
```

// JOB INTFOCOI
// % INIFGCO2
// * PROGRAM TD INITIALIZE STUCENT GRACE FILE ANC CLEAR STANCARD FILE. INTFGCC3
1/*
// DUP
*DELETE INTFG INTFGCOG
// FCR
*NAME INTFG
*ONE WORD INTEGERS
*EXTENCED PRECISION
*LIST ALL
*IOCS(DISK)
INTEGER ONE(160),7WC(40)
OEFINE FILE I (24,160,U,NXREC)
OEFINE FILE 2(8CO,40,U,NXRCC)
1.CO 10 I= 1,160
10 ONE(I)=0
OD 20 I=1,40
20THO(I)=0
NXREC=1
NXRCC = 1
TWO(1)=1
WRITE(Z'NXRCC) TWO
TWO(l) = 0
OD 30 1a1,24
30 WRITE(I'I) ONE
OD 40 I=2,800
40 WRITE(2'NXRCC) THC INTFGC28
CALL EXIT
END
/1 XEQ L Ol
*FILES(1,FSTDG),(2,SMSTU) INTFGC32
INTFGCO4
INTFGCOS
INTFGCC7
INTFGCCY
INTFGCIC
INTFGCIC
INTEG
INTFGCl2
INTFGCI3
INTFGC13
INTFGC14
INTFGCl5
INTFGC16
INTFGC17
INTFGC18
INTFGCl9
INTFGC19
INTFGCZC
INTFG021
INTFGC22
INTFGC23
INTFGO24
INTFGC25
INTFGC26
INTFGC27
INTFGC28
INTFGC29
INTFGC3C
INTFGC31

```
```

/1 JLE
1/ \$
// \# prGGrAN TC reAl a SET CT CAIA FCR TrE STUCENT PRCGRANS TC 'REAC'
1/ \& INTO a file ( lNE CF 12.)
// \#
// FLR
\#NAME INUFG

# IUCS(CARL,CISK,14C3 PRINTER)

\&EXTENCLU PRECISICN
*ONE hCRD INTEGERS
\#LISI SOURCE PROGRAN
*LISI SUBPRCGHAN NAHES
*LIST SYMBCLL TABLE
INTEGER BLFF(2),OELFF:`1
INTEGER A(216b), INPUT(160),NREM(77),CATA(212)
INTEGER DATAl(106),DATAZ(1C6)
INTEGER IA\&LE(16)
INIEGER ERR,EA
COMNUN A, INPUT, NREN,DATA
EQuIvalence (1able(1),A(2116))
EQUIVALENCE(LCCIL,Al1|),{LCC12,A(1001)}
EGUIVALENCE (NI,A(2114)),(NC,A(21151)
EQUIVALENCE (ERR,A(2109)),(EA,A(2025))
ECUIVALENCE (DATAI(L),DATA(1)),(DATAZ(I),CATA(IC7))
CEFIAE FILE 5(12,106,U,NXRDC)
1 NI=2
NO=b
REAC(NI,13) TABLE,NCTST
13 FORMAT(16A1,11)
CO E I=1,212
8 EATA (I) = C
EA=1
CO b I=1,106
2 CALL RDRGO
IF(ERR) 3,4,3.
3 PAUSE 7CO9
GO TC 2
4 ~ C A T A 1 ( 1 ) = L O C 1 1 ~
CATAE(I)=LOCI2
BUFF(1)= LCC11
BUFF(2) = LOC12
CALL EECEB(BUFF,OBUFF)
hRITE(NC,11) 1,DBUFF
ll FORHAT(1H ,13,3X,7A1)
5 CONTINUE
hRITE(5'2*NDTST-1) DATA
STOP 7777
END
// XEO L OL
\#FILES(5,SIMOT,OO15)
012345678G- ++\# 2

```

INDFGCOI
INDFGCO2
INDFGCC 3
INDFGCO4
INDFGCOS
INDFGCO6
I NDFGCO7
INOFGCOB
I NOFGCO9
INDFGG10
INDFGCII
INDFGCI2
INDFGC13
INDFGC 14
INDFGC 15
[NDFGO16
INDFGC17
INDFGC 18
INDFGC19
INDFGG20
1 NOFGO21
INDFGC22
INDFGC23
INDFG024
INDFG025
INDFGC26
INDFGC27
INDFG028
I NDFG029
I NDFGO 30
INDFGC31
INDFGG32
INDFGC 33
INDFGO34
INDFGO35
INDFGO36
INDFGC37
INDFGO38
INDFGC39
INDFGC40
INDFGO41
INDFGC42
INDFG043
INDFGO44
INDFGO45
INDFGO45
INDFGO46
INDFGC47
INDFGO48
INDFGC49
INDFGC49
INDFGO51



\begin{tabular}{|c|c|c|c|}
\hline \multicolumn{4}{|l|}{Pa, 1} \\
\hline // JOB & \multicolumn{2}{|l|}{CC26 0015} & 0015 \\
\hline log grive & CART SPEC & CART AVAIL & Priy Crive \\
\hline OOOC & 0026 & 0026 & 0001 \\
\hline 0001 & OC15 & 0015 & 000C \\
\hline \[
\begin{aligned}
& \text { / ASM } \\
& \text { \#LIST }
\end{aligned}
\] & & & \\
\hline
\end{tabular}



PAGE 4

01230165800122 012500 C 5800000 0127 O OCE5 012801 4C080130 012A O GCE7 012 B 014 C 300130 0120 C CODF 012E C. 90E2 012F 04838 013001610 01310184000001 013301800 01340 C007 0135 O1 OC000000
\(01370 \quad \mathrm{C} 802\)
\(013800440000 F 2\)
01340144000395
\(013 C 00\) E5800001
Ol3E 0 0000
\begin{tabular}{lll} 
& \\
\(013 F\) & 00 & \(C 5800001\) \\
0141 & 0 & \(00 C C\) \\
0142 & 0 & 7102 \\
0143 & 01 & \(600001 F 3\)
\end{tabular}

014500 C400000E
0147 0 DOCB
01480 90C8
C149 0104000240
014 B 0066800078


PAGE 5
\begin{tabular}{lll}
\(C 140\) & \(C\) & \(6 A 5 F\) \\
\(C 14 E\) & \(C\) & \(C 2 C 1\) \\
\(C 14 F\) & \(C\) & \(0 C C 4\) \\
0150 & \(C\) & \(C 2 C A\) \\
0151 & \(C 1\) & \(04 C 003 C 8\) \\
\(C 153\) & 01 & \(4 C 28016 A\)
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline 0155 & 0 & COBC \\
\hline 0156 & 0 & 9CBC \\
\hline C157 & c & 920日 \\
\hline 0158 & 01 & 0400036E \\
\hline C15A & 0 & C20A \\
\hline 0158 & 0 & 8C50 \\
\hline C15C & Cl & 9400036 E \\
\hline Cl5E & c & 0088 \\
\hline C15F & 01 & 8400036E \\
\hline Cl61 & 0 & \(8 \mathrm{C49}\) \\
\hline 0162 & 01 & 04000360 \\
\hline C164 & c & CCAE \\
\hline 0165 & 01 & 94000360 \\
\hline C167 & 0 & 9CAC \\
\hline 0168 & & DCAC \\
\hline 0169 & 0 & 7COF \\
\hline
\end{tabular}



AM 1415
AN 142 C AM 1425 AM \(143 C\) AM 1435 AM 1440 AM 1445 AM 1450 AM 1455 AM 146 C AM 1465 AM 1470 AM 1475 AM 148 C AM 1485 AM 1490 AM 1495 AM 1500 AM 1505 AM 1510 AM 1515 AM 1520 AM 1525 AM 1530 AM 1535 AM 1540 AM 1545 AM 1550 AM 1555 AM 1560 AM 1565 \(\Delta M 1570\) AM 1575 AM 1580 AM 1585 AM 1590 AM 1595 AM 1600 AM 1605 AM 1610 AM 1615 AM 1620 AN 1625 AM 1630 AM 1635 AM 1640 AM 1645 AM 1650 AM 1655 AM \(1 \in 6 C\) AM 1665 AM 1670 AM 1675 AM 168C AM 1685 AM 1690 AM 1695 AM 1700


ClaO CC C4COCCDO CIA2 C OC54
CIA3 C CC2C
ClA4 CC 04000000 Cla6 C C2l8
\(\begin{array}{lll}C 1 A 7 & C & \text { OC24 }\end{array}\)
OLAB O CC22
CIAS G DC18
CIAA \(0 \quad 7 C 13\)
\(\begin{array}{lll}\text { CIAB } & C & O C C G \\ \text { CIAC } & C & O C 7 A\end{array}\)
CIAD O OCOO

CIAE C ICOO
ClAF C 0818
CIBO CI 74C101C8
0182 C 3 CCO
CIB3 C OCCC
ClB4 C ot11
C185 C ECCC
C186 014 C200183
CIB8 O1 74FFOIC4
CIBA C 7CC5
CIBB \(C \quad 3 \mathrm{CCO}\)
ClBC C CCCO
CIBD C DCCA
CIBE C CCC4
CIBF C DCC4
OICO CI 4C4COLAE

\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|}
\hline & & & \multicolumn{5}{|l|}{} & AM & 1995 \\
\hline 01 C 2 & 0 & OCCO & H0COO & LC & & 10C00 & CONSTANT & \(A^{M}\) & 2000 \\
\hline 0163 & 0 & \(001 F\) & D31 & OC & & 31 & CONSTANT & AM & 2 CO 5 \\
\hline 0164 & 0 & 0 COO & PCNT & OC & & \#-* & PRINT CCUNT (NC OF CHAR.) & AM & 2C10 \\
\hline \(01 \mathrm{C6}\) & & 0000 & IOCC4 & BSS & E & 0 & IOCC TO SENSE LSW ANC & AM & 2C15 \\
\hline C1C6 & 0 & 0000 & & OC & & 0 & * RESET CSW ANC ILSh & AM & 2C20 \\
\hline 01 C 7 & 0 & OFCl & & OC & & /OFC1 & *FOR CONSOLE PRINTER. & AM & 2C25 \\
\hline OLC8 & & 0000 & IOCCP & BSS & E & C & IOCC TO PRINT CN CONSCLE & AM & 2030 \\
\hline 0168 & 1 & 0005 & & DC & & M8UF & *Printer one character at & AM & 2 C 35 \\
\hline 0169 & 0 & 0900 & & OC & & 10900 & *LCCATICN MBUF & AM & 2 C 40 \\
\hline 01 CA & 1 & 0005 & AMBUF & OC & & MBUF & ACORESS OF MELF & AM & 2045 \\
\hline 01 CB & 1 & 0183 & ANKC & OC & & NKC & ACORESS CF NKC & AM & 2050 \\
\hline O1CC & 0 & 0000 & SAVKC & OC & & *-* & LOCATION TC SAVE K/C ISS ENT & AM & \(2 C 55\) \\
\hline O1CE & & \(0 \mathrm{CO2}\) & I OAR2 & BSS & E & 2 & TOP OF CISK BUFFER 2 & AM & \(2 \mathrm{C6C}\) \\
\hline 0100 & 1 & 0102 & ANL. 5 & OC & & Ni 5 & ADORESS OF NEW LEVEL 5 & AM & 2065 \\
\hline \multirow[t]{2}{*}{0101} & 1 & 024E & AENT 5 & OC & & ENT5 & ACCRESS CF ENT 5 ENTRY PT & AM & 2C70 \\
\hline & & & \multicolumn{4}{|l|}{} &  & AM & 2075 \\
\hline 0102 & 0 & 9COO & NL5 & OC & & *-* & NEW ENTRY PCINT FOR LEVEL 5 & AM & 2080 \\
\hline 0103 & 0 & 0862 & & XIO & & IOCC & SENSE OEVICE STATUS hGRC L 5 & AM & 2085 \\
\hline 0104 & C & 1001 & & SLA & & 1 & SHIFT INT RUN BIT INTC ACC C & AM & 2090 \\
\hline 0105 & 02 & 4C100160 & & BSC & L & BOSCP, - & GO TO BCSCP IF NOT INT RUN & AM & 2095 \\
\hline 0107 & C & CCF9 & & LU & & AENT5 & LOAC ENTRY ACERESS FOR L 5 & AM & 2100 \\
\hline 0108 & 00 & 04000c00 & & STO & L & 13 & STORE INTC LCC 13 & AM & 2105 \\
\hline 010A & 0 & CCF1 & & LO & & SAVKC & LOAL SAVEC CON/KEYBC ISS ENT & AM & 2110 \\
\hline \multirow[t]{2}{*}{0108} & 0 & 0218 & & STO & X2 & 'ITCK & RESTGRE CON/KEYBO ISS TV ENT & AM & 2115 \\
\hline & & & \multicolumn{2}{|l|}{} & *** & ********* &  & AM & 2120 \\
\hline 0106 & 01 & 4400037 F & & BS I & L & IONC & WAIT FOR ALL I/C CFF & AM & 2125 \\
\hline cioe & Cl & C4000368 & & LC & L & CMI & loac ninus cne & AM & 2130 \\
\hline O1E0 & 01 & 94000114 & & S & L & LCCMM & SUBTRACT LEAGTH OF CCMMCN & AM & 2135 \\
\hline OlE2 & 0 & 0066 & & STO & & BCOMM +1 & Store as begining of conncin & AM & 2140 \\
\hline 01 E3 & Cl & 65800114 & & LOX & 11 & LCOMM & LOAC XRI WITH LEN. CF CCMMON & AM & 2145 \\
\hline Cle5 & C & 6287 & & LOX & 2 & -121 & ENTER XR2 WITH-121 & AM & 2150 \\
\hline 01 E6 & 02 & C6000078 & GAGN & LO & L2 & GTBL+12C & C LCAL value frcr buffer & AM & 2155 \\
\hline O1E8 & 00 & 05000000 & BCOMM & STO & L1 & *-* & StORE IN COMNCN & AM & 2160 \\
\hline OlEA & 0 & 7201 & & MDX & 2 & +1 & MOOIFY XR2 8Y 1, SKIP IF ZERC & AM & 2165 \\
\hline OlEB & 0 & \(7 \mathrm{CO2}\) & & MOX & & ARCUN & GO TC AROUN IIF NC SKIP) & AM & 2170 \\
\hline 01 EC & Cl & 740201EF & & MOX & L & INSCH, + 2 & 2 NOCIfy eranct accress ey +2 & AM & 2175 \\
\hline O1EE & 0 & 71FF & AROUN & MOX & 1 & -1 & MOCIFY XRI EY -1, SKIP If ZERC & AM & 2180 \\
\hline \multirow[t]{2}{*}{01 EF} & 0 & 70F6 & INSCH & MOX & & GAGN & GO TC GAGN (IF NC SKIP) & AM & 2185 \\
\hline & & & \multicolumn{2}{|l|}{**********} & * & ********* &  & AM & 2190 \\
\hline \(01 F 0\) & 0 & 4878 & & BCSC & & +-Z & SKIP ANC CFF INTERRUPT & AM & 2195 \\
\hline ClFl & C & 1000 & & NDP & & & NO-CP & AM & 2200 \\
\hline \multirow[t]{2}{*}{01F2} & 00 & 4C000000 & AMSR & BSC & L & *-* & EXIT FRCM ANS AND RETURN & AM & 2205 \\
\hline & & & ****** & *** & * & ********* &  & \(A^{\prime}\) & 2210 \\
\hline \(01 F 4\) & C & 0010 & H0010 & OC & & 10010 & CCNSTANT & AM & 2215 \\
\hline CIF5 & C & 8 CCC & H800C & OC & & 18000 & CONSTANT & AM & 2220 \\
\hline 01 F6 & 0 & 3 Ccc & H300C & CC & & 13000 & CCNSTANT (EGUALS haIt INSTR) & AM & 2225 \\
\hline CIF7 & 0 & 000C & SAVL 5 & CC & & *- & LOCATION TC SAVE L 5 TV & AM & 2230 \\
\hline ClF8 & C & 0000 & WAITC & OC & & 0 & WAIT IF NEGATIVE & AM & 2235 \\
\hline \multirow[t]{2}{*}{CIF9} & - & 000C & WAITO & OC & & 0 & WAIT INCICATCR FOR CON ENT Sh & AM & 224 C \\
\hline & & & \multirow[t]{2}{*}{*} & & & & haIt IF CON ENT Sh C IS LP & AM & 2245 \\
\hline O1FA & & OCOO & & BSS & \(E\) & 0 & EVEN CCRE ECLINCARY & AM & 2250 \\
\hline CIFA & 1 & 02F9 & \multirow[t]{2}{*}{ICCCD} & OC & & WAITD & REAC INTC WAITC & AN & 2255 \\
\hline \(01 F 8\) & C & 3AOC & & OC & & /3A00 & the conscle entry shitcres & AM & 226 C \\
\hline ClFC & C & 1 CCC & NOP & NCP & & & A NC-CP INSTRLCTICN & \(\Delta N\) & 2265 \\
\hline ClFO & C & OCOC & LAOOR & DC & & 0 & ACCRESS CF LAST INSTRUCTICN & AM & 2270 \\
\hline \multirow[t]{2}{*}{ClFE} & & OCC2 & LINST & BSS & \(E\) & 2 & LAST INSTRUCTICN & AM & 2275 \\
\hline & & & ****** &  & *** &  &  & AM & 2260 \\
\hline
\end{tabular}


0200 C 621 B C201 01 44C00395 0203 C \(5 C F C\) C204 Cl 4C2003AC C206 C C831 C 207 C 9834 0208 O1 4C1803AC C2CA C C\&20 C20B 0 882E C20C C D82B 0200018 C 6 C2OE C IECA C20F CC D4000C01 \(0211 \mathrm{ClC4COC2CI}\) \(02130105 C 00 C 43\)
C2l5 C CCEG C216. OC1L C217 C OBE2 0218 C C8ID C219 C EBCF C21B C1 4C100227 C21D Ci C400c2C1 C2LF C 1CC4 C2ZC C 1804 0221 C 4802 C222 C EEC2 0223 E E8C2
C224 C DCC3
C225 OL \(440 C 037 F\)
C 227 C 40 C 3
0228 C OCOC

C229 CI 4CCCO24E
C.2.2B C COOC





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\begin{tabular}{|c|c|c|}
\hline \multirow[t]{6}{*}{0294} & \multirow[t]{6}{*}{C} & \multirow[t]{6}{*}{DCAF} \\
\hline & & \\
\hline & & \\
\hline & & \\
\hline & & \\
\hline & & \\
\hline & & \\
\hline C295 & C & 6100 \\
\hline 0296 & C & CC2B \\
\hline C297 & 0 & ECB2 \\
\hline C298 & C & 4820 \\
\hline C299 & C & 611 C \\
\hline 029A & 0 & Cc27 \\
\hline C29B & 0 & ECAF \\
\hline C29C & 01 & 4C1802A3 \\
\hline C29E & C & 7110 \\
\hline C29F & 0 & Cc22 \\
\hline C2AO & 0 & ECA8 \\
\hline 02A1 & C & 4620 \\
\hline C2A2 & C & 71F8 \\
\hline C2A3 & 0 & COIE \\
\hline C2A4 & 0 & 1808 \\
\hline C2A5 & C & 1061 \\
\hline c2A6 & 00 & 04000002 \\
\hline 0248 & 01 & CE000084 \\
\hline C2AA & C & 1900 \\
\hline C2AB & C & 180C \\
\hline C2AC & C & 9 ClB \\
\hline C2AD & Cl & 4C2802C9 \\
\hline 02AF & C & 1801 \\
\hline C280 & C & 180F \\
\hline C2B2 & 0 & OC13 \\
\hline C2B2 & 0 & 18 C 3 \\
\hline C2B3 & C & 1800 \\
\hline C2B4 & OU & 04000c01 \\
\hline C2B6 & 01 & C5COC23F \\
\hline C2B8 & 0 & EC94 \\
\hline C289 & 0 & UCOC \\
\hline C2BA & 0 & \(4 \mathrm{Ca4}\) \\
\hline C2BB & Cl & \(4 C 000200\) \\
\hline C2BO & Cl & 4C0003AC \\
\hline
\end{tabular}


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PAGE 15
\begin{tabular}{lll} 
O2EB & 0 & CODA \\
O2EC & 01 & \(4 C 200365\)
\end{tabular}
\begin{tabular}{lll} 
O2EE & C & COD7 \\
C2EF & C & rC77
\end{tabular}
\begin{tabular}{lll} 
O2FO & 01 & \(4 C 180365\) \\
C2F2 & 6 & \(4 C 7 C\) \\
C2F3 & 0 & FC00 \\
O2F4 & 0 & \(440 C\) \\
C2F5 & 0 & \(7 C 29\) \\
\(02 F 6\) & 0 & 6106 \\
C2F7 & 0 & COCE \\
C2F8 & 01 & F50000CE
\end{tabular}
02FA Cl 4C180313


\begin{tabular}{|c|c|c|c|c|c|c|c|c|}
\hline C 342 & c & 7CLC & & nex & N503 & GC TC N5C3 If Check false & AM & 4605 \\
\hline C343 & Cl & C4CCC2C3 & & LC & L INST+1 & PUt ACORESS PCRTION INTC ACC & AM & 4610 \\
\hline C345 & c & 9 C 27 & & S & BCALL & Subtract accr beg. cf call & AM & 4615 \\
\hline C346 & Cl & 4C2803AC & & BSC & \(L \quad C G A,+z\) & go to cea on rinus & AM & 4620 \\
\hline C348 & 01 & 94000115 & & S & L LCLTV & SUBTRACT LENGTH OF CALL & AM & 4625 \\
\hline C34A & Cl & 4C1CC3AC & & BSC & L CGA,- & GC tc cga ch nct minus & AM & 4630 \\
\hline \multirow[t]{15}{*}{C34C} & \multirow[t]{15}{*}{c} & \multirow[t]{15}{*}{7 Cll} & & MOX & N508 & GC TC N508 & AM & 4635 \\
\hline & & & \multicolumn{4}{|l|}{***************************************************} & AM & 4640 \\
\hline & & & & Effe & Ctive accres & S IN LIEf T & AM & 4645 \\
\hline & & & \multicolumn{4}{|l|}{*****************************************************} & AM & 4650 \\
\hline & & & \multicolumn{4}{|l|}{* If the effective accress is within the} & AM & 4655 \\
\hline & & & \multirow[t]{2}{*}{L} & Iff TR & ANSFER VECTC & Cr, the instruction, index & AM & 4660 \\
\hline & & & & HREE, & ANC THE EFF & ective acoress are testec & AM & 4665 \\
\hline & & & T & C Cete & bmine if It & IS A PrCPER EATRY INTO the & AM & 4670 \\
\hline & & & & IbF TV & , if the te & est falls, It IS treatec as & AM & 4675 \\
\hline & & & L & F THE & EA has hith & In tree call transfer vector. & AM & 4680 \\
\hline & & & & \(F\) THE & test is succ & CESSFUL, then the mon & GM & 4685 \\
\hline & & & & NCICAT & OR IS SET TC & C IncICATE that it is valid & AM & 4690 \\
\hline & & & & OR the & PrLGRAM TC & be within the subroutine & AM & 4695 \\
\hline & & & \multirow[t]{2}{*}{\[
\begin{aligned}
& \text { * } 0 \\
& * * * *
\end{aligned}
\]} & R LIBF & TV AREAS. & & & 4700 \\
\hline & & & & ****** & ************ & ****************************** & AM & 4705 \\
\hline C340 & c & 4C21 & \[
\begin{aligned}
& * * * * * \\
& \text { N507 }
\end{aligned}
\] & \multirow[t]{2}{*}{BSI
CC} & BITS & RETURN ERANCH TO BITS & AM & 4710 \\
\hline 034 E & C & FFOO & \multirow[t]{3}{*}{} & & /FF00 & CHECK FOR SHORT BSI INSTR. & & 4715 \\
\hline C34F & c & 43 CC & & CC & 14300 & WITH XR3 (01000011). & AM & 4720 \\
\hline 0350 & 0 & \(7 \mathrm{Cl1}\) & \multicolumn{2}{|r|}{\multirow[t]{2}{*}{\(\operatorname{MDX}_{\text {L }}\)}} & N510 & GO TO N5IC If CHECK false & AM & 4725 \\
\hline C351 & 00 & C4000C03 & & & 3 & LOAC INCEX 3 & AM & 4730 \\
\hline 0353 & c & FC74 & \multicolumn{2}{|r|}{EOR} & SPXR3+1 & CCMPARE WITH Proper value xr 3 & & 4735 \\
\hline C354 & 01 & 4C2003AC & \multicolumn{2}{|r|}{BSC} & L CGA,Z & GO TO CG ACTUAL IF NOT ZERO & AM & 4740 \\
\hline 0356 & 0 & CC17 & \multicolumn{2}{|r|}{L0} & BLBTV & LOAC LOW ENC AODR LIBF TV & AM & 4745 \\
\hline C357 & C1 & 540002 C & \multicolumn{2}{|r|}{S} & EA & SUBTRACT EFFECTIVE ADORESS & AM & 4750 \\
\hline C359 & c & 1890 & \multicolumn{2}{|r|}{SRT} & 16 & Shift into a two word operano & AM & 4755 \\
\hline C35A & c & a 8 CE & \multicolumn{2}{|r|}{C} & 03 & divide ey trree & AM & 4760 \\
\hline 035B & 0 & 1800 & \multicolumn{2}{|r|}{RTE} & 16 & PLACE EXt INTG ACC & AM & 4765 \\
\hline 035 C & c1 & 4C2003AC & & BSC & CGA, 2 & go to ccmpute gr if Not zerd & AM & 4770 \\
\hline C35E & c & CCO9 & \multirow[t]{2}{*}{N508} & \multirow[t]{2}{*}{LO} & OM1 & LOAC ACC WITH MINUS ONE & AM & 4775 \\
\hline \multirow[t]{4}{*}{C35F} & \multirow[t]{4}{*}{cl} & 040002 C & & & L MON & Store into mon incicator & AM & 4780 \\
\hline & & & \multicolumn{3}{|l|}{\multirow[t]{2}{*}{*}} & MON = 1 WHILE IN MONITOR & AM & 4785 \\
\hline & & & & & & MON \(=0\) WHILE IN MAINLINE & AM & 4790 \\
\hline & & & \multicolumn{2}{|l|}{*} & & MEN \(=-1\) WHILE IN SUBRDUTINES & AM & 4795 \\
\hline \multirow[t]{10}{*}{0361} & \multirow[t]{10}{*}{0} & 7603 & \multicolumn{2}{|r|}{MOX} & xxx & GO TO XEO & AM & 4800 \\
\hline & & & \multicolumn{4}{|l|}{*****************************************************} & AM & 4805 \\
\hline & & & * & \multicolumn{2}{|l|}{Effective accre} & SS IN CALL T V & AM & 4810 \\
\hline & & & ***** & ****** & *********** & ******************************* & AM & 4815 \\
\hline & & & \multirow[t]{2}{*}{} & F tine & INSTRUCTION & IS OF A TYPE THAT ALTERS & AM & 4820 \\
\hline & & & & ORE, 1 & T hill Not & BE PERMITTEC TC EXECUTE. & AM & 4825 \\
\hline & & & * I & FIt 1 & S. NCT CF A & TYPE THAT ALTERS CORE, IT & AM & 4830 \\
\hline & & & \multirow[t]{2}{*}{*} & ILL BE & PERMITTEO & to execute, without an entry & AM & 4835 \\
\hline & & & & N THE & garbage tab & LE. * & AM & 4840 \\
\hline & & & \[
{ }_{* * * *}^{*}
\] & ****** & ************ & ****************************** & AM & 4845 \\
\hline \multirow[t]{2}{*}{C362
\(C 364\)} & \multirow[t]{4}{*}{01
0} & 740002 C 5 & \multirow[t]{2}{*}{N510} & \multirow[t]{2}{*}{Mcx
Mox} & \(L\) Store,o & SKIP IF Store indic. Is zero & AM & 4850 \\
\hline & & 7047 & & & CGA & go to cga if NCT LERD & AM & 4855 \\
\hline \multirow{2}{*}{C364} & & & \multicolumn{2}{|l|}{* MOX} & & Store incicatcr = O If loac & AM & 4860 \\
\hline & & & \multicolumn{2}{|l|}{\multirow[b]{2}{*}{XXX BSC}} & & Store incicatcr \(=1 \cdot\) IF Store & AM & 4865 \\
\hline \multirow[t]{4}{*}{0365} & \multirow[t]{4}{*}{01} & 4C80020F & & & 1 TSTEA & ExIt frcm test ea rcutine & AM & 4870 \\
\hline & & & \multicolumn{4}{|l|}{****************************************************} & AM & 4875 \\
\hline & & & \multicolumn{4}{|l|}{\multirow[t]{2}{*}{\begin{tabular}{l}
* constants for use by lower half \\
*****************************************************
\end{tabular}}} & AM & 4880 \\
\hline & & & & & & & AM & 4885 \\
\hline 0367 & 0 & \(0 \subset 32\) & IOCT & CC & \$IOC T & & AM & 4890 \\
\hline
\end{tabular}

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\begin{tabular}{lll}
0368 & 0 & FFFF \\
0369 & 0 & \(00 C 3\) \\
\(036 A\) & \(C\) & \(O C 40\) \\
\(C 36 B\) & \(C\) & \(O C 32\) \\
\(036 C\) & 0 & \(O C O C\) \\
\(C 360\) & 0 & \(O C O C\) \\
\(036 E\) & 0 & 0000
\end{tabular}
\begin{tabular}{lll}
\(036 F\) & 0 & \(0 C 00\) \\
0370 & 01 & \(C 40002 C 2\) \\
0372 & 0 & 6909 \\
0373 & 01 & \(6580036 F\) \\
0375 & 0 & \(E 100\) \\
0376 & 0 & \(F 1 C 1\) \\
0377 & 0 & 4818 \\
0378 & 0 & 7101 \\
\(G 379\) & 0 & 7102 \\
\(037 A\) & 0 & 6903 \\
0378 & 00 & 65000000 \\
0370 & 00 & \(4 C 000000\)
\end{tabular}


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\begin{tabular}{lll}
\(C 345\) & \(C\) & \(0 C G C\) \\
\(C 396\) & \(C\) & \(08 C 3\) \\
\(C 397\) & \(C\) & \(C C C 4\) \\
\(C 398\) & 01 & \(4 C 80 C 395\) \\
\(C 39 A\) & & \(0 C C C\) \\
\(C 39 A\) & 1 & \(C 1 S C\) \\
\(C 39 B\) & \(C\) & \(3 A C C\) \\
\(C 39 C\) & \(C\) & \(O C C C\)
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline C390 & 0 & 0 CCO \\
\hline C39E & C & 4 CFE \\
\hline C39F & C & ECCR \\
\hline C3AC & Cl & 4し58C390 \\
\hline C3A2 & C & CCCb \\
\hline C3A3 & C & COC3 \\
\hline C3A4 & & \\
\hline C349 & Cl & 4 CBCO 39 D \\
\hline C348 & C & OCC2 \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|}
\hline PAGE & 20 & & & & & & & & \\
\hline C3C5 & 0 & 6B2F & & Stix & 3 & XR3 & STORE XR3 INTC XR3 & AM & 5475 \\
\hline 03 C 6 & 0 & 6A2D & & STX & 2 & XR2 & STORE XR2 INTC XR2 & AM & 5480 \\
\hline 03 C 7 & 00 & 67000000 & SPXR3 & LOX & 13 & *-* & RESTGRE PRCPER VALUE XR3 & AM & 5485 \\
\hline 03C9 & 0 & CB7A & & LOD & 3 & 122 & LOAC FIRST 2 WCRCS GF FAC & AM & 5490 \\
\hline 03CA & 0 & 0823 & & Sto & & SAVF1 & SAVE FDR CUTPUT ROUTINE & AM & 5495 \\
\hline 03CB & 0 & CB7C & & LOD & 3 & 124 & LGAC SECGNC 2 WCRES CF FAC & AM & 5500 \\
\hline 03CC & 0 & 0823 & & STO & & SAVF2 & SAVE FDR CUTPUT RCUTINE & AM & 5505 \\
\hline 03 CO & 0 & CB7E & & LOO & 3 & 126 & LOAC THIRC 2 HORCS DF FAC & AM & 5510 \\
\hline O3CE & 0 & 0823 & & STO & & SAVF3 & SAVE FDR DUTPUT ROUTINE & AM & 5515 \\
\hline 03CF & 01 & CC000238 & & LOO & \(L\) & INSCT & LDAE LCUBLE INSTRUCTION CT & AM & 5520 \\
\hline 0301 & 0 & A833 & & 0 & & ClE4 & OIVICE ey lccec & AM & 5525 \\
\hline 0302 & 01 & DC000238 & & STO & \(L\) & INSCT & Store ocuble instr count & AM & 5530 \\
\hline 0304 & 00 & C4000032 & & LO & \(L\) & \$IOCT & LDAC I/C BUSY INCICATOR & AM & 5535 \\
\hline 0306 & 0 & DCIF & & STC & & SIGCT & SAVE IN SICCT & AM & 5540 \\
\hline 0307 & 0 & 1010 & & SLA & & 16 & CLEAR ACC & \(A M\) & 5545 \\
\hline 0308 & OC & 04000032 & & STO & L & \$IDCT & CLEAR I/O BUSY INCICATOR & AM & 5550 \\
\hline 030A & 00 & 04C000EE & & STO & \(L\) & \$OBSY & CLEAR CISK BuSY INCICATCR & AM & 5555 \\
\hline 030C & 0 & 6106 & & LEX & 1 & 6 & ENTER INCEX 1 hITH E & AM & 556 C \\
\hline 0300 & 0 & C824 & OLOOP & LOD & & ICAR & LDAC SECTOR LENGTH AND ACCR. & AM & 5565 \\
\hline 030E & 0 & CCIA & & LO & & C320 & LCAC 32C AS SECTOR LENGTH & AM & 557 C \\
\hline 03DF. & 01 & D08003FB & & STO & 11 & CPAR-1 & STORE AT TCP CF BUFFER & AM & 5575 \\
\hline 03E1 & 01 & CC0003FA & & LCO & Ll & OPAR-2 & LCAC CISK Paraneters & AM & 5580 \\
\hline 03E3 & 00 & 440000 F2 & & BSI & L & D2000 & GO TC CISK RCUTINE & AM & 5585 \\
\hline C3E5 & 01 & 74010403 & & HOX & \(L\) & IOAR+1,1 & MOCIFY SECTCR AOCRESS BY ONE & AM & 5590 \\
\hline 03E7 & 0 & \(71 F E\) & & MOX & 1 & -2 & MCDIFY XRI BY -2, SKIP IF ZERC & AM & 5595 \\
\hline \(03 \mathrm{E8}\) & 0 & \(70 F 4\) & & MOX & & OLCOP & GC TC CLCOP CN NO SKIP & AM & 5600 \\
\hline C3E9 & 00 & 040A41E3 & & LINK & & O8UG T & CALL LINK TO CBUGT & AM & 5605 \\
\hline O3EE & & 0002 & SAVF1 & 8SS & E & 2 & LOCATIDN TC SAVE FAC & AM & 5610 \\
\hline 03F0 & & \(0 \mathrm{CO2}\) & SAVF2 & BSS & E & 2 & LCCATION TC SAVE FAC & AM & 5615 \\
\hline 03F2 & & 00 C 2 & SAVF3 & BSS & E & 2 & LCCATION TC SAVE FAC & AM & 5620 \\
\hline 03F4 & 0 & 0 COO & XR2 & OC & & *-* & location to save grade inc. & AM & 5625 \\
\hline 03F5 & 0 & OCOC & XR3 & DC & & *-* & location tc save enc value & AM & 5630 \\
\hline & & & * & & & & *DF INCEX 3 & AM & 5635 \\
\hline 03F6 & 0 & OCOC & SIOCT & OC & & *-* & LCCATION TC SAVE I/C BUSY INE & AM & 564 C \\
\hline C3F7 & 0 & EEEE & heeee & CC & & /EEEE & CONSTANT & AM & 5645 \\
\hline 03F8 & 0 & OC20 & H2O & DC & & 120 & CONSTANT & AM & 5650 \\
\hline 03F9 & 0 & 0140 & C320 & CC & & 320 & CCNSTANT & AM & 5655 \\
\hline 03FA & 0 & 2000 & H2000 & LC & & 12000 & CONSTANT & AM & 5660 \\
\hline 03FC & & 0000 & H0001 & BSS & E & 0 & CONSTANT (ONE) & AM & 5665 \\
\hline 03FC & & 0000 & OPAR & BSS & E & 0 & TABLE DF CiSk parameters & AM & 5670 \\
\hline C3FC & 0 & 00C1 & & OC & & 1 & CISK hrite & AM & 5675 \\
\hline 03FD & 1 & 031 C & & DC & & ICAR 3 & ALORESS CF CISK BUFFER 3 & AM & 568 C \\
\hline 03FE & 0 & 0001 & & CC & & 1 & CISK hrite & AM & 5685 \\
\hline 03FF & 1 & O1CE & & LC & & IOAR2 & ACORESS CF CISK BLFFER 2 & AM & 5690 \\
\hline 0400 & 0 & 0061 & & CC & & 1 & CISK hRITE & AM & 5695 \\
\hline 0401 & 1 & 0000 & & CC & & IGARI & ADCRESS CF CISK BUFFER l & AM & 5700 \\
\hline 0402 & 31 & 22065109 & \[
{ }_{*}^{\text {IOAR }}
\] & DSA & & SAVGR & \begin{tabular}{l}
CISK FILE LENGTH, SECTOR \\
*aCCRESS, ANC AC CF SECTCRS.
\end{tabular} & AM & \[
\begin{aligned}
& 5705 \\
& 571 C
\end{aligned}
\] \\
\hline 0405 & 0 & 2710 & 0154 & DC & & 10000 & CONSTANT & AM & 5715 \\
\hline COOL & & & -CMCN & EQU & & 1 & LENGTH CF COMMCN & AM & 5720 \\
\hline C004 & & & 'HWET & EGU & & 4 & LENGTH CF CCRE IMAGE HEACER & AM & 5725 \\
\hline 0008 & & & - TVWC & EGU & & 8 & LENGTH CF TRANSFER VECTCR & AM & 5730 \\
\hline C009 & & & 'hCNT & EGU & & 9 & LENGTt CF CCRE LCAC & AM & 5735 \\
\hline COOA & & & - XR3X & EGU & & 14 & SETTING FOR INCEX 3 & AM & . 5740 \\
\hline 0018 & & & - ITCK & EGU & & 24 & INTERRUPT ENTRY TC K EC/CCN PR & AM & 5745 \\
\hline OCOE & & & \$CORE & EGU & & /000E & SIZE CF CCRE & AN & 575 C \\
\hline 0028 & & & \$PRET & EQU & & 10028 & PRE-CP I/G ERRCR TRAP & AM & 5755 \\
\hline C032 & & & \$IOCT & EGU & & 10032 & I/C BUSY INCICATOR & \(\Delta M\) & 576 C \\
\hline
\end{tabular}
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